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Collateralized Networks

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Abstract. This paper studies the spread of losses and defaults in financial networks with two interrelated features: collateral requirements and alternative contract termination rules. When collateral is committed to a firm’s counterparties, a solvent firm may default if it lacks sufficient liquid assets to meet its payment obligations. Collateral requirements can, thus, increase defaults and payment shortfalls. Moreover, one firm may benefit from the failure of another if the failure frees collateral committed by the surviving firm, giving it additional resources to make other payments. Contract termination at default may also improve the ability of other firms to meet their obligations through access to collateral. As a consequence of these features, the timing of payments and collateral liquidation must be carefully specified to establish the existence of payments that clear the network. Using this framework, we show that dedicated collateral may lead to more defaults than pooled collateral, we study the consequences of illiquid collateral for the spread of losses through fire sales, we compare networks with and without selective contract termination, and we analyze the impact of alternative resolution and bankruptcy stay rules that limit the seizure of collateral at default. Under an upper bound on derivatives leverage, full termination reduces payment shortfalls compared with selective termination.

Keywords: contagion • OTC markets • financial regulation • network • fire sales • collateral • automatic stays for qualified financial contracts

1. Introduction

This paper studies the spread of losses and defaults through financial networks when payment obligations are at least partly secured by collateral. A combination of changes in regulation and industry practices following the financial crisis of 2007–2009 have greatly expanded the use of collateral in trading and lending. In the over-the-counter (OTC) derivatives market, most standardized contracts now trade through central counterparties, which require participants to post collateral in the form of initial margin (IM). The part of the market that continues to trade bilaterally is now also subject to IM requirements. Similarly, unsecured interbank lending is far lower than its precrisis levels and has mostly been replaced by collateralized lending through repurchase agreements. See, for example, Duffie (2017), Ghamami and Glasserman (2017), Financial Stability Board (2017), and U.S. Department of the Treasury (2017b) for background.

Collateral provides a buffer against the spread of losses: if one party to a contract defaults on a payment obligation, its counterparty can seize available collateral to offset the loss. In this sense, collateral supports financial stability.

But, as our analysis shows, this is not the whole story. We highlight additional mechanisms that complicate the impact of collateral. First, committing collateral to specific contracts and counterparties can lead to an ex post inefficient allocation of a firm's assets: a firm may find itself unable to make a current payment obligation on one contract despite having posted collateral to protect future potential obligations on other contracts. Firms do not ordinarily have the option to terminate contracts to recover posted collateral. In reducing counterparty credit risk, collateral requirements can increase strains on funding liquidity because collateral requirements create additional funding needs. These effects are intimately connected to contract termination rights because terminating a contract can provide access to collateral and, indeed, is ordinarily necessary for access to collateral.

Collateral held in less liquid assets creates a further consequence for contagion. At the failure of one institution, its counterparties would seize and liquidate collateral. This sell-off or fire sale could drive down the price of the collateral assets, creating market losses at other firms holding those assets. In particular, firms that had posted similar assets as collateral would find themselves with a collateral shortfall—and
an obligation to add collateral—as a result of the fire sale. See Stein (2013), Shleifer and Vishny (2011), and Gennaioli and Shleifer (2018, chapters 1 and 2), and the references therein for background on the role of fire sales in financial crises.

We develop these ideas in a network model. The nodes of the network are parties to financial contracts; for brevity, we sometimes refer to these as banks though we have in mind a broader set of financial and even nonfinancial companies. The nodes are linked through contracts that carry payment obligations. We take the network configuration as an input to our analysis; our model does not seek to explain how a particular configuration comes about.

We build on the standard framework of Eisenberg and Noe (2001). The Eisenberg–Noe model takes a set of nodes with balance sheets linked through uncollateralized payment obligations and identifies one or more clearing vectors. A clearing vector describes a set of actual payments under which a node never pays more than it owes, all contracts have equal seniority, and nodes face limited liability. These properties lead to a fixed-point characterization of clearing payments. The Eisenberg–Noe model has been extended to cover many other features, including bankruptcy costs (Rogers and Veraart 2013) and claims of different seniority (Elsinger 2009) or maturity (Kusnetsov and Veraart 2019); see Cabrales et al. (2016), Glasserman and Young (2016), Hurd (2016), and Jackson and Pennock (2019) for overviews and extensive references. Despite these many extensions, the inclusion of collateral poses special complications and requires a departure from the usual solution approach.

In a collateralized network, the failure of one node may improve the ability of other nodes to meet their obligations. If a surviving node had committed collateral to a contract with the failed node, the failure frees that collateral, providing the surviving node additional resources to make other payments. Indeed, the freed collateral might even be necessary for the surviving node to meet its obligations, in which case the failure of one node prevents the failure of another. We show that, under these circumstances, the notion of clearing payments may not be well defined.

In an equilibrium of the Eisenberg–Noe model, the timing of events (payments and defaults) is immaterial, and these events may be understood as occurring simultaneously. In modeling collateral, we separate the timing of two types of events following a default. We assume that creditors have immediate access to collateral posted by the defaulting node, but the freeing of collateral posted to the default node follows a short delay. This modeling choice is supported by key principle 5 of the Basel Committee on Banking Supervision and International Organization of Securities Commissions (2015, p. 20) principles on margin requirements for noncleared derivatives: “Initial margin collected should be held in such a way as to ensure that (i) the margin collected is immediately available to the collecting party in the event of the counterparty’s default, and (ii) the collected margin must be subject to arrangements that protect the posting party to the extent possible under applicable law in the event that the collecting party enters bankruptcy.” The return of collateral in (ii) does not carry the same urgency as the access to collateral in (i), so we do not assume they occur simultaneously.

By separating these events, we show that we can arrive at a well-defined set of clearing payments. The separation eliminates the possibility that the failure of one node could prevent the failure of another node by freeing collateral—a scenario we see as unrealistic as well as a complication for the analysis. Along with the clearing payments, we characterize the set of defaulting nodes and the redistribution of collateral.

We use this framework to evaluate the effects on a network of various policy options related to collateral and contract termination rights because contract termination controls access to collateral. We take as our measures of financial stability the size of the default set and the network’s total payment shortfall. We examine these measures in various scenarios and establish the following conclusions:

i. Costless termination. As a starting point for comparison, we consider networks in which nodes are free to terminate contracts, and we show that, under this assumption, posting collateral is equivalent to making certain payments, so collateral plays no essential role. This reference point establishes the close connection between access to collateral and restrictions on contract termination that runs through the rest of our analysis.

ii. Pooled collateral. We investigate the trade-off between committing collateral to specific counterparties versus pooling collateral and holding additional cash. This trade-off is analogous to the comparison between collateral requirements and capital requirements: capital absorbs any type of loss, whereas collateral protects specific obligations. Although pooling may appear to allow a better allocation of resources, we show that pooling is guaranteed to reduce defaults under additional conditions. For instance, we show that, when committed collateral exceeds current payment obligations, pooling reduces defaults and does not affect payment shortfalls. This result is applicable with derivative contracts. Because, in OTC derivatives markets, collateral in the form of initial margin captures part of extreme potential future exposures, it can exceed current payments (Ghamami 2020a). In general, however, pooling may produce larger or smaller payment shortfalls, depending on a node’s position in
the network, so the trade-off cannot be resolved by considering a node in isolation.

iii. Collateral illiquidity. We expand our model to capture the potential spread of losses through collateral fire sales. We extend the method of Cifuentes et al. (2005), which models illiquidity through a price-impact function. When a node fails, its creditors liquidate collateral, driving down its price and lowering the value of similar assets held by other nodes. This price-mediated channel amplifies losses beyond the direct effect of missed payments. Collateral illiquidity increases defaults and payment shortfalls.

iv. Automatic stays. We consider the effect of a stay under which payments are made before collateral is liquidated and show that it has no effect on defaults or payment shortfalls when collateral is liquid. We interpret this point as consistent with the policy recommendations of Duffie and Skeel (2012) and the subsequent finalized stay rule on collateral sale in repo markets, under which collateral can be accessed immediately only if it is held in cash or cash-like assets.

v. Accelerated payments triggered by defaults. Shortly after Lehman Brothers filed for bankruptcy, its counterparties terminated approximately 733,000 of more than 900,000 OTC derivatives contracts (Fleming and Sarkar 2013). A contract termination creates a new payment obligation, equal to the market value of the contract, from the out-of-the-money party to the in-the-money party. Lehman’s counterparties generally terminated contracts with positive value to the surviving party and chose not to terminate contracts with positive value to Lehman. We expand our model to compare this type of selective termination (ST; known as “cherry-picking”) with full termination (FT). Incorporating accelerated payments through contract termination complicates the analysis, again because one node can benefit from the failure of another, in this case because a default by one node may accelerate payments to other nodes. Arriving at a well-defined set of clearing payments requires making further assumptions on the timing of events; we assume a delay between the failure of a node and any accelerated payments resulting from that failure. This modeling choice is supported by the unfolding of the Lehman bankruptcy in 2008. Claims against Lehman resulting from contract termination became part of the bankruptcy process, leading to delays in payments as discussed in Fleming and Sarkar (2013).

vi. Alternative contract termination rights. Using the framework of accelerated payments, we compare networks under different contract-termination protocols and discuss these comparisons in the context of postcrisis automatic stay rules on access to collateral upon a counterparty’s failure. We compare three scenarios: selective termination (cherry-picking) by surviving nodes, full termination of all contracts with a failed node, and no termination. We show that selective termination ordinarily results in fewer defaults, but we also show that full termination can reduce systemwide payment shortfalls under a constraint on a measure of firms’ derivatives leverage. We make a similar comparison of full-termination and no-termination scenarios. These comparisons are motivated by continuing discussions on the treatment of derivatives in the bankruptcy and resolution of large financial institutions. See Jackson (2012), Ghamami (2020b), U.S. Department of the Treasury (2018), and the references therein.

Section 2 reviews the Eisenberg–Noe model, explains the difficulties introduced by collateral, and presents our solution; it also examines the case of costless contract termination and collateral pooling. Section 3 introduces illiquid collateral and presents the joint solution of clearing payments and market prices for collateral assets along with the implications for defaults and payment shortfalls. Section 4 extends our model to cover accelerated payments from contract termination. Section 5 compares defaults and payment shortfalls under alternative termination scenarios governing collateral access. We defer all proofs to the appendix.

2. Network Model

In this section, we first review the model of Eisenberg and Noe (2001) and then introduce collateral.

2.1. Networks Without Collateral

We consider a network with nodes $N = \{1, \ldots, N\}$ representing banks or other market participants. (We can also think of one node as representing the outside world.) We use the following notation:

\[ p_{ij} = \text{payment due from } i \text{ to } j, i, j \in N; \]
\[ c_i = \text{cash held by node } i \in N; \]
\[ p_{ij} = \text{actual payment from } i \text{ to } j, i, j \in N. \]

We refer to $c_i$ as cash for brevity; more generally, $c_i$ represents the near-term cash value of assets (other than the $p_{ij}$) available to node $i$ to make payments. We assume that payment obligations are netted so that $\bar{p}_{ij}$ and $\bar{p}_{ij}$ cannot both be strictly positive.

We seek to model default triggered by illiquidity rather than insolvency, and this differentiates our formulation from most interbank network models. In our setting, the claims and obligations $\bar{p}_{ij}$ and cash values $c_i$ are not intended to provide a complete accounting of a node’s balance sheet; each node would typically have other longer-term assets and liabilities. The $p_{ij}$ measure payments due. A node defaults if it does not have the cash to make a payment even if the total value of its assets exceeds the value of its
liabilities. Each $c_i$ includes the cash a node could raise by borrowing against long-term assets.

We imagine that, outside the model, some nodes have experienced an exogenous loss of asset value, and we proceed to evaluate payments made, taking the $c_i$ as cash levels after the exogenous shock. Given a collection of payments $p_{ij}$, node $i$’s cash is given by

$$A_i^0 = c_i + \sum_{k \neq i} p_{ik},$$

and its payments due are given by

$$L_i = \sum_{k \neq i} p_{ki}.$$  

(1)  

(2)

Node $i$ defaults if $A_i^0 < L_i$, so the default set is

$$D = \{ i : N : A_i^0 < L_i \}.$$  

(3)

If node $i$ defaults, its creditors’ claims all have equal priority, and any remaining cash held by node $i$ is paid to the creditors in proportion to their claims. These proportions are given by

$$d_{ij} = \frac{p_{ij}}{\sum_{k \neq i} p_{ik}}, \quad i, j \in N.$$  

(4)

For each $i$, we have $\sum_{j \neq i} d_{ij} = 1$.

Clearing payments are characterized by the fixed-point equation

$$p_{ij} = \bar{p}_{ij} \wedge \left[ c_i + d_{ij} \sum_{k \neq i} p_{ik} \right].$$  

(5)

This specification ensures that actual payments never exceed obligations ($p_{ij} \leq \bar{p}_{ij}$), all creditors have equal priority in the sense that they receive payments proportional to their claims in the event of default, and the total payments made $\sum_i p_{ij}$ cannot exceed the available cash $A_i^0(p)$. Eisenberg and Noe (2001) show the existence of a solution to (4) and also give conditions for uniqueness. Tarski’s (1955) fixed-point theorem ensures the existence of a largest and a smallest solution of (4).

2.2. Networks with Collateral: Round 1

We now introduce collateral, which we also refer to as IM. We let

$$m_{ij} = \text{margin posted by node } i \text{ to node } j.$$  

Suppose, for example, that nodes $i$ and $j$ are two banks that have entered into a swap contract. Under rules adopted in the United States in 2015 and 2016, each bank is required to post IM as collateral against potential future payments to the other bank. If node $j$ is a central counterparty (CCP), then $i$ posts IM to $j$, but $j$ does not post IM to $i$. Because IM is intended to cover potential future losses, a node often faces a margin requirement on a contract even if no payment is due.

We refer to $m_{ij}$ as margin posted or committed by node $i$ to node $j$. We assume that the margin $m_{ij}$ remains an asset of node $i$ until node $i$ fails to make a payment to node $j$. The quantity $m_{ij}$ differs from other assets held by node $i$ in that node $j$’s claim to the $m_{ij}$ has priority over the claims of any other creditors. We take $m_{ij} = 0$ for all $i$.

We assume the following sequence of events at default. Node $i$ goes into default when it has insufficient cash (including payments received from other nodes) to meet its obligations. Rather than make partial payments that would not stave off default, node $i$ briefly suspends making any payments. At this point, node $j$ seizes enough of the collateral $m_{ij}$ to cover any payment due $p_{ij}$ from $i$ to $j$. The amount of collateral seized by node $j$ from node $i$ is given by

$$\Delta_{ij} = \begin{cases} 
    m_{ij} \wedge \bar{p}_{ij}, & i \in D; \\
    0, & i \notin D.
\end{cases}$$  

In particular, if node $i$ defaults, node $j$’s claim to the collateral $m_{ij}$ is determined solely by the payment obligation $\bar{p}_{ij}$ and is unaffected by any other claims on node $i$’s assets. This is the defining feature of collateral.

If the margin seized by node $j$ from node $i$ is insufficient to cover the obligation $\bar{p}_{ij}$, node $j$ retains a residual claim of $\bar{p}_{ij} - m_{ij}$, which has equal priority with any residual claims against $i$ by other nodes. This claim produces a partial payment from $i$ to $j$ if node $i$ has any remaining assets. To reflect a pro rata allocation of node $i$’s cash to these equal-priority claims, we replace (3) with the proportions

$$d_{ij}^{(1)} = \frac{[\bar{p}_{ij} - m_{ij}] \wedge \sum_{k \neq i} [p_{ik} - m_{ik}]}{\sum_{k \neq i} [p_{ik} - m_{ik}]}.$$  

(6)

If the denominator is zero, no node has a residual claim on $i$, and we may set $d_{ij}^{(1)} = 0$; if node $i$ does not default, set $d_{ij}^{(1)} = d_{ij}^{(0)}$ as in (3).

Let $p_{ij}^{(1)}$ denote the total payment made by node $i$ to node $j$, consisting of any collateral seizure $\Delta_{ij}$ and any partial payments based on the proportions in (6). Given payments $p_{ij}^{(1)}$ made to node $i$, the cash available to node $i$ is given by

$$A_i^{(1)} = c_i + \sum_k p_{ik}^{(1)}.$$  

(7)

This expression has the same form as (1), but it includes node $i$’s access to collateral $m_{ij}$ posted by other nodes and seized according to (5). The set of defaulting nodes is given by

$$D = \{ i \in N : A_i^{(1)} < L_i \}.$$  

(8)
with the sets $L_i$ as in (2). The requirements for first-round clearing payments $p^{(1)}_{ij}$ now take the form

$$p^{(1)}_{ij} = \begin{cases} p_{ij} \wedge \left[ m_{ij} + \Delta_{ij}^{(1)} \right], & i \in D; \\ p_{ij}, & i \notin D. \end{cases}$$

(9)

These conditions extend (4) using (6)–(8) and including the collateral $m_{ij}$. If, for some $i$, $m_{ij} + \Delta_{ij}^{(1)} < p_{ij}$, then $\Delta_{ij}^{(1)} < p_{ij}$ and node $i$ is in default. We may, therefore, write (9) as

$$p^{(1)}_{ij} = \bar{p}_{ij} \wedge \left[ m_{ij} + \Delta_{ij}^{(1)} \right].$$

(10)

This expression is similar to the Eisenberg–Noe Equation (4), but the similarity hides an important difference: we cannot recover the default set from payments in (10) because a node in default may meet its payment obligations through access to collateral. In (4), $p_{ij} = \bar{p}_{ij}$ implies that node $i$ had not defaulted, but we cannot make a similar inference from $p^{(1)}_{ij} = \bar{p}_{ij}$ in (10).

We confirm the existence of clearing payments $p^{(1)} = \{p^{(1)}_{ij}, i \in \mathcal{N}\}$ satisfying (6)–(10) shortly, but we first elaborate on the potential complications introduced by collateral.

### 2.3. Freed Collateral: Round 2

Following a default by node $i$, it may happen that the collateral committed to node $j$ exceeds the payment obligation to $j$, in which case the excess $m_{ij} - \Delta_{ij}$ would become available to node $i$ to meet other obligations. If another node $k$ also defaults, then margin $m_{ik}$ posted by $i$ to $k$ may also become available to $i$. This return of collateral is potentially problematic for the existence of clearing payments as the following example illustrates.

**Example 2.1.** Consider the network illustrated in Figure 1 with nodes A, B, and C. The number inside each circle indicates the node’s cash; all cash levels are initially zero. The arrows indicate the directions of payment obligations, and the labels on the arrows indicate the amounts due. The labels in square brackets show posted margin. For example, the label “[5-A]” above node C indicates that node A has posted collateral worth five to node C. As no node has cash, all nodes default. The default of node C returns collateral worth five to A, and similarly, B and C each recover collateral worth five from the defaults A and B.

The result is the configuration in the middle of the figure. With the freed collateral, all nodes can meet their payment obligations, resulting in the final configuration. However, if all nodes have met their payment obligations, have they defaulted? It is not possible in this example to simultaneously determine a consistent set of payments and default designsations. In the rightmost configuration, no node appears to be in default, but we cannot reach that configuration without freeing collateral, which requires defaults.

The problem illustrated by this example is that the default of one node can lead to the return of collateral to other nodes. Those nodes may then be able to make greater payments to the defaulted node, potentially lifting that node out of default and precluding the existence of an internally consistent set of payments.

We avoid this difficulty through the definition of the default set in (8) based on $p^{(1)}$. In this formulation, freed collateral does not become available until after payments have been made and nodes have been declared in default. We have, thus, separated the payments process into two rounds: a first round in which payments are made and collateral is seized as needed and a second round in which excess collateral and collateral posted to defaulted nodes becomes available to the posting party to make other payments.

**Figure 1. Three-Node Network**

<table>
<thead>
<tr>
<th></th>
<th>A</th>
<th>B</th>
<th>C</th>
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<tr>
<td>[6]</td>
<td>S</td>
<td>S</td>
<td>S</td>
</tr>
<tr>
<td>[8]</td>
<td>S</td>
<td>S</td>
<td>S</td>
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<tr>
<td></td>
<td>A</td>
<td>B</td>
<td>C</td>
</tr>
</tbody>
</table>
| ![Notes. All nodes start with zero cash. The labels in brackets show collateral: the amount posted and the posting party.](image)

Obstacles to defining clearing payments arise in other extensions of Eisenberg and Noe (2001) as well. Elsinger (2009) redefines clearing vectors to cover cross-holdings of debt and equity among banks, David and Lehar (2017) consider clearing payments when debt is subject to renegotiation, Kusnetsov and Veraart (2019) propose a detailed algorithm to handle debt with different maturities, Jackson and Pernoud (2019) note that clearing payments may not be well defined when banks buy credit protection on other banks, and Banerjee and Feinstein (2018) similarly preclude banks from speculating on other banks. None of these extensions fits our setting. Bichuch and Feinstein (2019) consider networks with collateral, but in their setting, banks post collateral outside the network to raise cash; in particular, they do not allow one bank to seize collateral posted by another bank, which is a key feature of collateral in our setting.

In Figure 1, no collateral is seized because collateral is sitting at the wrong nodes to secure payments due. For example, C owes A, but A has posted collateral to C. As a result, all collateral is freed when the nodes default. Figure 2 shows the opposite configuration with the direction of collateral postings reversed. All nodes default, all nodes seize collateral worth five to cover the amount owed, and there are no remaining...
payment obligations. As in Figure 1, there is no way to simultaneously define clearing payments and the default set: if all nodes default, then all payments are covered, so no node defaults; if no node defaults, then no collateral is seized, no payments are covered, so all nodes default.

The return or freeing of collateral in Figure 1 implicitly assumes that contracts with defaulted nodes are automatically terminated. Recall that margin \( m_{ij} \) is intended to cover future potential losses on a contract (e.g., a swap), so it is possible to have margin posted when no payment is due; that is, \( m_{ij} > 0 \) and \( p_{ij} = p_{ij} = 0 \). If node \( k \) defaults on its payments to other nodes, the implications for node \( i \) are unclear. We assume (for now) that any contracts between \( i \) and \( k \) are canceled and that posted collateral is returned: \( m_{ik} \) becomes available to node \( i \), and \( m_{ki} \) becomes available to node \( k \). We examine alternative assumptions on contract termination and collateral access later.

Suppose the first round results in payments \( p_{ij}^{(1)}, i, j \in N \). In other words, the \( p_{ij}^{(1)} \) satisfy (6)–(9). If the payment due from \( i \) to \( j \), \( p_{ij} \), exceeds the payment \( p_{ij}^{(1)} \), then node \( i \) enters the second round with a remaining payment obligation to node \( j \):

\[
p_{ij}^{(2)} = p_{ij} - p_{ij}^{(1)}.
\]  

(11)

In round 2, we deal with the allocation of remaining resources to meet remaining payment obligations. These remaining resources result from freed collateral.

In the Eisenberg–Noe setting, (4) ensures that a node in default pays out all its cash. But the use of collateral in round 1 makes two types of resources potentially available at a defaulted node in round 2: a node may recover excess collateral committed to another node, and it may recover freed collateral committed to a defaulting node as a result of contract termination. Figure 3 illustrates these mechanisms.

**Example 2.2.** Figure 3 shows three starting configurations. In the first network, A defaults, B seizes collateral to cover missed payments from A, and excess collateral is returned to A, allowing A to pay C. In the second network, C defaults on its obligation to B; C’s default cancels its contract with A, freeing the collateral posted by A to C, leaving A with five in cash. In the last network, A defaults, B seizes its collateral, and A is left with four units of cash despite having defaulted. See Section 2.4 for a discussion of full repayment following default.

The total collateral “returned” to node \( i \) after round 1 is given by

\[
r_i = \frac{\sum_{j \in D} (m_{ij} - \Delta_{ij})}{\sum_{j \notin \emptyset} m_{ij}}, \quad i \in D, \quad i \notin D.
\]  

Here, \( D \) is fixed by (8) in round 1. If this freed collateral is insufficient to meet all of node \( i \)'s remaining claims, payments are made in the proportions

\[
a_{ij}^{(2)} = \frac{a_{ij}^{(2)}}{\sum_{k \in D} p_{ik}^{(2)}},
\]  

(13)

taking \( a_{ij}^{(2)} = 0 \) if the denominator is zero. Thus, we seek two payments satisfying

\[
p_{ij}^{(2)} = p_{ij}^{(2)} \land a_{ij}^{(2)} \left( r_i + \sum_{k \in D} p_{ik}^{(2)} \right), \quad i, j \in N.
\]  

(14)

For later use, we record a relationship between the allocation fractions in the two rounds.

**Lemma 2.1.** If \( a_{ij}^{(2)} \neq 0 \), then \( a_{ij}^{(2)} = a_{ij}^{(2)} \).

In the following, we use \( p_{ij}^{( \ell )} = \sum_{\ell} p_{ij}^{( \ell )}, \ell = 1, 2 \), to denote the total amount paid by node \( i \) in each round.

**Proposition 2.1.** For any collateral levels \( m_{ij}, i, j \in N \), the two clearing payments \( (p_{ij}^{(1)}, p_{ij}^{(2)}) \) satisfying (6)–(9) and (11)–(14). Moreover,

\[
p_{ij}^{(1)} + p_{ij}^{(2)} = p_{ij} + A \left( p_{ij}^{(1)} + p_{ij}^{(2)} \right) + \sum_{k \in D} m_{ik},
\]  

(15)

where

\[
A(p_{ij}^{(1)} + p_{ij}^{(2)}) = c_i + \sum_{k \in D} \left( p_{ki}^{(1)} + p_{ki}^{(2)} \right).
\]

As our examples illustrate, the existence of clearing payments depends on separating the timing of payments due to the freeing of collateral: the default set \( D \) and returned collateral \( r_i \) are determined by the first-round payments \( p_{ij}^{(1)} \). Without this separation, the network often fails to admit a consistent set
of clearing payments. Equation (15) has a simple interpretation: the total payments made by a node over two rounds equal the lesser of the node's total obligations and the node's total cash, including collateral. However, (15) does not extend to node-specific payments $p_{ij}^{(3)} + p_{ji}^{(3)}$ because collateral is initially committed to specific counterparties.

In subsequent sections, we compare outcomes of networks under different policies regarding collateral and contracts. We make these comparisons based on the sets of defaulting nodes and the system-wide payment shortfall, which builds on the total payments in (15).

**Definition 2.1.** A network's payment shortfall is the difference between payments due and payments made given by

$$L = \sum_i \left( p_i - p_i^{(1)} - p_i^{(2)} \right) = \sum_{i,j} \left( p_{ij} - p_{ij}^{(1)} - p_{ji}^{(2)} \right).$$

(16)

For each node $i$ that defaults, $p_i - p_i^{(1)} - p_i^{(2)}$ is the difference between $i$'s total payment obligation and its total payments; if $i$ does not default, then $p_i = p_i^{(1)}$, $p_i^{(2)} = 0$, and its shortfall is zero.

Table 1 shows the default sets and payment shortfalls for the examples in several figures. In several cases, we have designed the examples to have $L = 0$ to highlight the effect of two rounds of payments. In the first example of Figure 3, for instance, increasing node $A$'s payment obligations from 9 to 10 + $x$ results in $L = x$ for any $x \geq 0$ without changing $D$.

To see that our two rounds cannot be combined in general, consider an ordinary Eisenberg–Noe network without collateral. Suppose that, for some node $i$, there exists a cash level $\xi$ such that $i$ defaults if $c_i < \xi$ and $i$ does not default if $c_i \geq \xi$ (We can see from (7) and (8) that such thresholds commonly exist.) Now, introduce collateral and consider any proposed protocol with the following two intended features: (i) collateral posted by a node is used toward the node's payment obligations and if only if that node defaults, and (ii) a node is deemed to default if and only if it fails to meet its payment obligations. Such a protocol is not in general consistent.

**Proposition 2.2.** Suppose $\sum_j m_{ij} > \xi - \xi_i > 0$ and suppose $m_{ik} = 0$ for all $k \neq i$. If node $i$ defaults, then it meets all its payment obligations; if node $i$ does not default, then it does not meet all its payment obligations. In other words, no choice of default set $D$ is consistent with the protocol.

The proof of this claim is simple: If node $i$ is deemed to default, then the resources $\xi_i + \sum_j m_{ij} > \xi_i$ suffice to meet the node's obligations along with any cash received from other nodes. If node $i$ does not default, then the cash level $\xi_i < \xi_i$ does not suffice for $i$ to make its payments. In either case, we have a contradiction. Similar contradictions result from many other configurations.

2.4. Default with Full Repayment

We have seen (as in Table 1) that it is possible for a node to default in the first round yet fully meet its payment obligations by the end of the second round, eliminating its first-round shortfall. In other words, a node may fail through illiquidity—a shortage of cash to meet payments due—even if it is solvent because some of its assets have been pledged as collateral. Indeed, the major failures and near-failures of 2008 are generally understood as (at least initially) crises of liquidity rather than solvency. Post-crisis regulation responded by introducing a liquidity coverage ratio for banks as a complement to traditional capital requirements.

One may ask whether a default is costly if creditors are ultimately repaid. This question goes to our definition of the default set, so we highlight two important practical considerations that inform our modeling choice. First, a delay in payments can be costly because of its ripple effect on downstream parties that rely on receiving those payments to meet their own obligations. Delays can also lead to credit downgrades, which then affect a node's ability to borrow. Second, a delay that results in bankruptcy destroys franchise value; it locks a failed firm out of markets that rely on a solid reputation for meeting payment obligations on time, and such a reputation is not restored simply through eventual repayment. Nor are other costs of bankruptcy recovered. The specifics of these mechanisms are beyond the scope of our model, but these considerations explain why we treat first-round defaults as costly even when second-round obligations are fully met.

2.5. Networks with a Free Termination Option

As a benchmark, we consider a variant of our model in which nodes have the option to terminate contracts

<p>| Table 1. Default Sets $D$ and Payment Shortfalls $L$ for Examples in the Figures |
|-----------------------------|-----------------------------|-----------------------------|-----------------------------|-----------------------------|</p>
<table>
<thead>
<tr>
<th>Figure 1</th>
<th>Figure 2</th>
<th>Figure 3</th>
<th>Figure 5</th>
<th>Figure 6</th>
</tr>
</thead>
<tbody>
<tr>
<td>$D = {A, B, C}$</td>
<td>$D = {A, B, C}$</td>
<td>$D = {A}$</td>
<td>$D = {A}$</td>
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<tr>
<td>$L = 0$</td>
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<td>$L = 5$</td>
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Note. Figure 3 has three examples.
in order to redeploy collateral they have posted. This termination option is exercised whenever a node would otherwise default. We investigate other contract termination scenarios more extensively in Sections 4 and 5.

To motivate this variant, consider the three networks in Figure 3. Suppose that node A has the option to terminate any contract by using collateral posted to pay its counterparty, recovering any remaining collateral. In the first case, this would mean paying five to B and recovering five; in the second case, paying zero to C and recovering five; and in the third case, paying five to B and recovering zero. In each case, A would avoid default.

We show in Appendix A.3 that this model can be reduced to a standard Eisenberg–Noe model with lower payment obligations and adjusted cash balances in the following sense:

**Lemma 2.2.** Payments \( p_{ij} \) clear the network with collateral and free contract termination if and only if payments \( p_{ij} = (m_{ij}, A, p_{ij}) \) clear the reduced Eisenberg–Noe network. The two networks have the same default sets and payment shortfalls.

The reduced network is defined precisely in Appendix A.3. This result shows that collateral plays no essential role in a network in which each node can recover collateral by terminating contracts. With the option to terminate, posting collateral reduces to paying down certain obligations. This result allows us to compare the collateralized networks of Sections 2.2 and 2.3 with otherwise identical networks that allow free contract termination.

**Proposition 2.3.** Free contract termination reduces defaults but does not affect total payment shortfalls.

This comparison confirms the natural intuition that collateral trapped in the “wrong” places increases defaults (because it is not immediately available to meet payment obligations) but does not affect eventual payments (because excess collateral is eventually returned and deployed to make payments).

The free-termination model is useful for illustration, but it overlooks two key points. First, parties to swaps and similar contracts do not ordinarily have the right to terminate contracts unilaterally and must pay for that option by, for example, buying a swaption. The issue of termination near default can be particularly contentious—a point we return to in Section 5. Furthermore, in the second example of Figure 3, suppose we change A’s obligation to B to eight and change C’s obligation to B to three. Node A would like to terminate its contract with C to recover its collateral. But C has no reason to agree and would like to extract a payment from A. To avoid default, both nodes would like to keep at least three of the five units of collateral.

The second key point missed by this model variant is the distinction between payments due and contract values. Suppose node A has taken out a loan of 50 from node C and posted five in collateral. In the middle example of Figure 3, no interest is due on the loan. However, if A were to terminate the contract to recover its collateral, it would need to repay the principal, dramatically increasing its immediate payment obligations. The same can happen with a cross-currency swap, which entails an exchange of principal at maturity. We consider these consequences of contract termination in detail in Section 4. For these reasons, we work with the framework of Sections 2.2 and 2.3 for the rest of the paper: collateral becomes accessible only following a default.

### 2.6. Collateral Pooling

In the examples of Figures 2 and 3, some defaults occur because collateral is tied up in the wrong places. More precisely, some defaults could be avoided if nodes were able to hold on to their collateral as cash rather than commit it to specific counterparties. This leads to a trade-off similar to the trade-off between capital requirements and collateral requirements: collateral provides a buffer against specific losses, whereas cash absorbs any type of loss. Similar considerations apply in debates over “ring-fencing” capital to absorb losses in specific jurisdictions as opposed to holding capital at the parent level.

These considerations might suggest that defaults and losses can be reduced by having all nodes hold additional cash rather than post collateral. But that conclusion is incorrect because changing the distribution of collateral changes the distribution of payments.

Consider the example of Figure 4. The labels on the edges indicate payment obligations. None of the nodes holds cash, but node A may have posted collateral to \( B_1 \) or \( C_1 \). We compare defaults and shortfalls in the following scenarios:

- \( c_A = 0, m_{AB_1} = 0, m_{AC_1} = 2 \) : \( D = \{A, B_1, B_2\}, L = 6 \);
- \( c_A = 2, m_{AB_1} = 0, m_{AC_1} = 0 \) : \( D = \{A, B_1, B_2, C_1\}, L = 5 \);
- \( c_A = 0, m_{AB_1} = 2, m_{AC_1} = 0 \) : \( D = \{A, C_1\}, L = 4 \).

**Figure 4.** Pooling Collateral May Increase or Decrease Payment Shortfalls and the Number of Defaults

![Figure 4](image-url)
In the first scenario, collateral is committed to node $C_1$. Holding the collateral as cash instead, as indicated in the second scenario, results in a payment of one to each of $B_1$ and $C_1$. This increases the number of defaults, but it reduces systemic payment shortfalls, which argues in favor of pooling. However, the third configuration, with collateral committed to node $B_1$, yields the fewest defaults and the smallest shortfall.

As this example illustrates, pooling collateral is not unambiguously better or worse than committing it to specific counterparties. The comparison depends on the network and cannot be resolved by considering a node in isolation.

The next result shows that pooling is preferable in two settings. To be precise, we need some terminology. Let us say that one network is obtained from another by pooling collateral if it results from one or more transformations of the form $m'_{ij} = m_{ij} - \delta$, $c'_i = c_i + \delta$, $\delta \leq m_{ij}$. By the excess collateral posted by node $i$ to node $j$, we mean $[m_{ij} - p_{ij}]^+$. We say that the network has proportional collateral if $m_{ij} = k_i p_{ij}$ for some $k_i \in [0,1]$ for all $i$ and $j$. (The case $k_i > 1$ would be a special case of excess collateral.)

**Proposition 2.4.** (i) Pooling excess collateral reduces defaults and does not affect payment shortfalls. (ii) Under proportional collateral, pooling that preserves proportionality reduces defaults and does not affect payment shortfalls.

This proposition and the example of Figure 4 together suggest that pooling is unambiguously better (in reducing defaults and shortfalls) only under special conditions. A proportional collateral rule (such as a fixed loan-to-value ratio) is found in some circumstances, but it is not applicable with derivative contracts that carry different levels of risk or when payment obligations are not known precisely at the time collateral is posted. Excess collateral is applicable to derivative contracts as collateral in the form of initial margin can exceed current contractual payment obligations. Because initial margin in part captures extreme potential future exposures, it can exceed current payments, making the excess collateral condition particularly relevant to OTC derivatives markets.

### 3. Illiquid Collateral and Fire Sales

In Sections 2.2 and 2.3, we implicitly treat collateral as cash: if node $j$ seizes collateral $\Delta_{ij}$ from node $i$, node $i$'s payment obligation is reduced by exactly $\Delta_{ij}$. In the noncleared derivatives market, a wide range of less liquid securities, including corporate bonds, foreign-denominated bonds, and equities, are accepted as collateral, and these types of securities are also used as collateral in repurchase agreements.

In this section, we extend our earlier analysis to incorporate the use of less liquid collateral. When a creditor seizes collateral, it must sell the collateral to recover cash. Selling less liquid collateral drives down its price, spreading losses to other holders of similar assets. Moreover, if excess collateral is available, the creditor has no incentive to sell at the best possible price, making the risk of a fire sale particularly acute.

Indeed, collateral liquidity is at the heart of debates over contract termination rights and bankruptcy stays that motivate our investigation. As explained, for example, in Roe and Adams (2015), the purpose of bankruptcy stays is to avoid the value destruction and fire sales that would occur if creditors were allowed to seize and sell a failing firm’s illiquid collateral. Regulators and industry participants continue to debate the extent to which less liquid collateral should be allowed for derivatives and repo and whether the nature of the collateral necessitates different restrictions on contract termination and rules on stays.

#### 3.1. Round 1 Revisited

As a first step, we reformulate the analysis of Sections 2.2 and 2.3 to incorporate illiquidity. We now take $m_{ij}$ to be the shares of collateral committed by node $i$ to node $j$ for an asset with price $\pi_i$, making $m_{ij}/\pi_i$ the value of the collateral. The case considered in Sections 2.2 and 2.3 corresponds to a constant value $\pi = 1$. In this section, the price starts at one but falls as collateral is liquidated.

We posit that $\pi$ is a strictly decreasing function $G(1, \Lambda)$ of the total shares sold, $\Lambda$, the first argument of $G$ indicating the initial price of one. To be concrete, we set

$$\pi = G(1, \Lambda) \equiv e^{-a\Lambda},$$

for some $a > 0$. A larger $a$ corresponds to a less liquid asset. This choice of price-impact function is also used in Cifuentes et al. (2005); Amini et al. (2016) discuss conditions on $a$.

We make the simplifying assumption that all collateral is held in a single illiquid asset. Assigning different prices and price-impact functions to different collateral assets would complicate notation without significantly changing our analysis. Assuming a single illiquid asset for all collateral oversimplifies the effect of fire sales, but the overstatement can be offset through a smaller value of $a$. The choice of $a$ should reflect the average price impact across different types of collateral and the imperfect correlation in price impact across different securities.

As in (7) and (8), we have

$$A_{ij}^{\pi} = c_i + \sum_k p_{kij}^\pi, \quad D = \left\{ j : A_{ij}^{\pi} < \sum_k p_k \right\},$$

but the payments in $A_{ij}^{\pi}$ received by node $i$ now reflect the market price of the collateral asset. If node $i$ defaults and node $j$ seizes collateral $m_{ij}$ with price $\pi_j$,
node $j$ now holds a residual claim of $(\bar{p}_j - \pi m_{ij})^*$ against node $i$, reflecting the market value $\pi m_{ij}$ of the collateral available. Node $i$'s assets are allocated to other nodes in proportion to the values of these claims, so we replace (6) with

$$a_i^{(1)} = \frac{(\bar{p}_j - \pi m_{ij})^*}{\sum_{k \neq i} (\bar{p}_{ik} - \pi m_{ik})^*}. \quad (19)$$

**Example 3.1.** Figure 5 illustrates the difference between allocation proportions in (6) and (19). Node $A$ has six cash in and 20 in payment obligations, so it defaults. Node $B$ seizes the five shares of the collateral asset posted by $A$. The six in cash held by $A$ is divided between $B$ and $C$. Under (6), $B$ would be allocated a proportion $(10 - 5)/15 = 1/3$ and $C$ a proportion 2/3. Under (19), node $B$ claims a proportion $(10 - 5\pi)/(20 - 5\pi)$, which is a decreasing function of the price $\pi$ at which $B$ sells collateral it seized from $A$; the allocation proportions depend on the market price of the collateral asset.

With asset price $\pi > 0$, the shares of collateral seized and sold by node $j$ upon the default of node $i$ are given by

$$\Delta_{ij} = \begin{cases} m_{ij} \wedge \bar{p}_j; & i \in D; \\ 0, & i \notin D, \end{cases} \quad (20)$$

which reduces to (5) with $\pi = 1$. Dividing by $\pi$ in (20) converts the dollar obligation $\bar{p}_j$ into the number of shares required to cover the payment at the current market price. For $\pi = 0$, interpret (20) as $\Delta_{ij} = m_{ij}$ if $\bar{p}_j > 0$ and $\Delta_{ij} = 0$ otherwise. The total shares of collateral liquidated are given by

$$\Delta = \sum_i \sum_j \Delta_{ij}, \quad (21)$$

and the sale of these shares drives down the price through (17). We modify (9) by seeking clearing payments $p^{(3)}$ and an asset price $\pi^{(3)}$ satisfying

$$p^{(3)}_j = \begin{cases} \bar{p}_j \wedge \pi^{(3)} m_{ij} + a_i^{(1)} A_i^{(1)}; & i \in D; \\ \bar{p}_{ij}, & i \notin D, \end{cases} \quad (22)$$

together with (17)–(21).

### 3.2. Round 2 Revisited

Suppose that round 1 clears with payments $p^{(1)}$ and asset price $\pi^{(1)}$. Define $r_i$ as in (12) and interpret it as the number of shares of collateral freed or returned to node $i$. As in (11), node $i$'s remaining obligation to node $j$ is given by $p^{(2)}_{ij} = \bar{p}_{ij} - r_i^{(1)}$. To meet its remaining obligations $\bar{p}^{(2)}_j$, node $i$ liquidates some or all of its freed collateral, which further drives down the price of the collateral. Given second-round payments $p^{(2)}_j$ and a second-round share price $\pi^{(2)} > 0$, the number of shares liquidated by node $i$ is given by

$$\Gamma_i = \tau_i \wedge \frac{1}{\pi^{(2)}} \left( \sum_{j \neq i} p^{(2)}_{ij} - \sum_{j \neq i} \pi^{(2)}_j \right). \quad (23)$$

The expression in parentheses is the difference between node $i$'s remaining payment obligations and the second-round payments it receives; dividing by $\pi^{(2)}$ yields the number of shares of collateral required to make up this shortfall, but node $i$ cannot liquidate more than the $\tau_i$ shares it recovers. The total amount liquidated by all nodes is $\Gamma = \sum_i \Gamma_i$, driving the price to

$$\pi^{(2)} = G(\pi^{(1)}, \Gamma) = \pi^{(1)} e^{-\alpha \Gamma}. \quad (24)$$

To clear round 2, we need payments $p^{(2)}$ and a price $\pi^{(2)}$ satisfying

$$p^{(2)}_j = \bar{p}_j \wedge \pi^{(2)} \sum_{k \neq i} p^{(2)}_{kj} + \pi^{(2)} \tau_i, \quad (25)$$

together with (23) and (24) and $a_i^{(2)}$ as in (13). The following result ensures the existence of clearing payments and compares networks with liquid and illiquid collateral.

**Proposition 3.1.** There exist clearing payments and prices $(p^{(1)}, \pi^{(1)})$ and $(p^{(2)}, \pi^{(2)})$ for rounds 1 and 2 with illiquid collateral.

Two features in particular distinguish this result from other network models with fire sales, such as Cifuentes et al. (2005), Brauwezel and Wagaleth (2018), and Cont and Scharanning (2017): one is the need to split the payments into two rounds because of the collateral, and the second is the fact that the proportions (19) depend on the collateral price $\pi$. Both features lead to more involved arguments for the existence of clearing payments and prices. The source of the fire sale is also different. In prior work, the fire sale is driven by banks selling their own assets to meet capital requirements; in our setting, the fire sale is driven by creditors selling collateral to recover payment shortfalls.

Our next result compares networks with liquid and illiquid collateral. In stating the result, we need to
account for the possibility that each network admits multiple sets of clearing payments. We show that each network has a largest and smallest set of first-round and total payments \((p^{(1)}, p^{(2)})\). The following comparison should be understood to hold for the largest and smallest solutions of networks with liquid and illiquid collateral:

**Proposition 3.2.** Collateral illiquidity increases defaults and the total payment shortfall.

### 3.3. Collateral Fire Sale and Contagion

Our model formulation in Sections 2.2 and 3.1 assumes that, upon a default in round 1, collateral is seized first, and partial payments are made second. For comparison, in this section, we consider an alternative formulation in which access to collateral is delayed and partial payments are made first. Within round 1, we reverse the order of collateral seizure and partial payments; round 2 proceeds as before. This reversal may be interpreted as the result of an automatic stay in which a defaulting node’s counterparties are prevented from immediately seizing and liquidating collateral. (We discuss stays in greater detail in Section 5.) We give an example here and leave the details for Appendix A.7. We show there that, in the absence of fire sales, the total payments made from one node to another remain unchanged under this protocol even though the mix of collateral and cash payments may change. With illiquid collateral, delaying collateral liquidation reduces systemwide losses and defaults.

Figure 6 illustrates the result. Node A defaults because its cash level eight falls below its payment obligations of 20. Under the collateral-first protocol, nodes B and C seize and liquidate the collateral posted by A, which is \(m_{AB} = 5\) and \(m_{AC} = 10\). Applying (10) with \(d_{ij}^{(0)} = 1\) and \(d_{ij}^{(1)} = 0\), node A makes a cash payment of five to node B in addition to the collateral transfer of five. The total first-round payments are \(p_{AB}^{(1)} = p_{AC}^{(1)} = 10\), and there are no second-round payments. Under the payments-first protocol, node A first makes a cash payment of eight to node B, again because \(d_{ij}^{(0)} = 1\) and \(d_{ij}^{(1)} = 0\). Node C liquidates \(m_{AC} = 10\) shares of collateral, but node B liquidates only two shares and returns three to node A. The total first-round payments are \(p_{AB}^{(1)} = p_{AC}^{(1)} = 10\), and there are no second-round payments. However, the total amount of collateral liquidated has been reduced from 15 to 12.

As in this example, the analysis of Appendix A.7 shows that, when collateral is held in cash-like assets, the spread of losses through the network is unaffected by the order of collateral seizure and partial payments. However, as the example suggests, the collateral-first protocol results in greater collateral liquidation. As a consequence, the collateral-first protocol can result in greater losses when collateral is illiquid.

If we interpret the payments-first protocol as the result of an automatic stay on collateral seizure by the counterparties to a failed node, then this observation is in line with policy recommendations of Duffie and Skeel (2012) and the subsequent finalized stay rule on collateral sale in repo markets. The stay rule allows immediate seize and liquidation of collateral only if it is held in cash or cash-like assets.

### 4. Accelerated Payment Obligations from Contract Termination

A financial firm’s failure to make a payment due on one contract may trigger the termination of other contracts on which no payments are due. This is particularly true in over-the-counter derivatives markets. OTC derivative contracts often provide participants the right to terminate a contract if the counterparty enters bankruptcy even if the counterparty has met all obligations under the contract. Bankruptcy courts also provide failed firms certain rights to terminate contracts. Upon termination, the market value of a swap or other derivative contract becomes due from the out-of-the-money party to the in-the-money party; in this sense, a default can accelerate payment obligations that would not otherwise be due.

In this section, we augment the model of Section 3 to incorporate this feature. We show that accelerated payments from contract termination can create inconsistencies in a network model similar to those we encountered with collateral. We again resolve these complications by carefully specifying the timing of events. Whereas clearing a network with collateral could require two rounds, addressing contract termination may require as many rounds as there are nodes because each round of terminations may trigger further defaults and, thus, further terminations. In this section, we assume that all contracts with a node are terminated upon the node’s default; in the next section, we address selective termination.

Let

\[ v_{ij} = \text{positive value to node } j \text{ of its derivative contracts with node } i. \]
Under full contract termination, the default of node \( i \) triggers the termination of its derivatives; at termination, node \( i \) incurs an obligation to pay node \( j \) the outstanding value \( v_{ij} \) if \( v_{ij} > 0 \). If, however, \( v_{ij} > 0 \), then the default of \( i \) triggers a payment obligation of \( v_{ij} \) from \( j \) to \( i \). We assume that all contracts between \( i \) and \( j \) are fully netted, so \( v_{ij} = 0 \).

Whereas the \( p_{ij} \) represent payments due under ordinary circumstances (including, for example, routine payments on swap contracts), the \( v_{ij} \) represent asset (\( v_{ij} > 0 \)) or liability (\( v_{ij} < 0 \)) values for node \( i \) that turn into payment obligations only upon contract termination. If all derivatives are subject to daily settlement (such as futures contracts), then all changes in market values would be offset by daily payments, and we would always have \( v_{ij} = 0 \). In practice, the balance sheets of large banks show significant derivatives assets and liabilities, indicating that not all contracts are settled daily.

The contingent payment obligations created by contract termination create some of the same complications we saw previously in the following example illustrates:

**Example 4.1.** Consider a three-node network with

\[
p_{12} = 4, \quad p_{13} = 2, \quad p_{23} = 2, \quad c_1 = c_2 = 4.
\]

All other model parameters are zero except for the termination values, which we specify shortly. Node 1 defaults as its cash, \( c_1 = 4 \), falls below its required payments to nodes 2 and 3. Node 1 pays \( p_{12} = 4 \) to node 2 and \( p_{13} = 2 \) to node 3. Node 3 defaults as the payment of \( p_{23} = 2 \) is more than \( p_{13} = 2 \). Prior to contract termination, node 2 does not have any payment obligations.

Suppose \( v_{21} = 2 \) and \( v_{13} = v_{23} = 0 \). Node 2, with cash of \( c_2 + p_{21} + p_{23} = 8 \), pays \( p_{21} = 2 \) to node 1. With this influx of cash, node 1 can now fully meet its payment obligations to nodes 2 and 3. In other words, node 1 can make all payments due, apparently avoiding default, but only if it defaults! This example shows that it is not always possible to simultaneously specify a consistent set of payments and default designations with automatic contract termination even without collateral.

To resolve this type of inconsistency, we separate payments into rounds as we did before, consistent with the sequence of events described informally in Example 4.1. However, each round of contract terminations can now potentially trigger additional payment obligations and, therefore, additional defaults. In a network with \( N \) nodes, we may have up to \( N \) rounds of defaults and \( N \) rounds of payments.

As in Section 3.3, we include the sale of any collateral in the total payment from node \( i \) to node \( j \) and denote this total payment by \( p_{ij} \). We assume the network follows the collateral-first protocol of Sections 3.1 and 3.2. Perhaps most importantly, we assume that, if node \( i \) defaults in round \( t \), then any termination values \( v_{ij} \) or \( v_{ij} \) triggered by this default become payment obligations in round \( t + 1 \). This timing is consistent with the interpretation of accelerated obligations as consequences of default rather than causes of default.

### 4.1. Round 1

Round 1 proceeds exactly as in Section 3.1. The first-round quantities \( p_{ij}^{(1)}, v_{ij}^{(1)}, A_i^{(1)}, D_i, a_i^{(1)}, \) and \( \Lambda^{(1)} \equiv \Lambda \) are defined by Equations (17)–(22).

### 4.2. Subsequent Rounds

We now consider round \( m, 2 \leq m \leq N \). With \( c_i^{(1)} = c_i \), the cash available to node \( i \) at the beginning of the round is given by

\[
c_i^{(m)} = c_i^{(m-1)} + \sum_{k \neq i} p_{ik}^{(m-1)} - \sum_{k \neq i} p_{ki}^{(m-1)},
\]

with \( p_{ij}^{(m)} = p_{ij} \). We let \( D_m \) denote the set of nodes that default in round \( m \) (and not before) and define \( S_m \) to be the set of nodes that survive rounds \( 1, \ldots, m \),

\[
S_m = \{ i : i \notin \bigcup_{l=1}^m D_l \}.
\]

with \( 1 \leq n \leq N \), and \( S_0 = \emptyset \). Payment obligations in round \( m \) are defined by

\[
p_{ij}^{(m)} = \begin{cases} v_{ij} + p_{ij}^{(m-1)} - p_{ij}^{(m-1)}, & i \text{ or } j \in D^{(m-1)}, \\ p_{ij}^{(m-1)} - p_{ij}^{(m-1)}, & i \text{ or } j \notin S_{m-1}; \\ 0, & i, j \in S_{m-1}. \end{cases}
\]

In each round, any previous payment obligation \( p_{ij}^{(m-1)} \) is reduced by any payment made \( p_{ij}^{(m-1)} \). A default by either node in the previous round creates the additional obligation \( v_{ij} \) from contract termination. If both nodes have survived to the current round, then the original payment obligation \( p_{ij} \) was met in the first round, and no subsequent obligation has been introduced, so the remaining obligation is zero.

As in Section 2.3, defaults in the previous round may free collateral in the current round. Recovered shares of collateral in round \( m \) are given by

\[
r^{(m)}_i = \frac{\sum_{i \in S_m \setminus i \in D^{(m-1)}} m_i - \Delta^{(m-1)}_i}{\sum_{j \in S_m \setminus j \in D^{(m-1)}} m_j}, \quad i \in D^{(m-1)}; \quad i \in S_{m-1}.
\]

The number of shares of collateral posted by \( i \) that are seized and liquidated by \( j \) in round \( m \) is given by

\[
\Delta^{(m)}_j = \begin{cases} m_j \wedge \frac{r^{(m)}_i}{\chi^{(m)}}, & i \in D^{(m)} \text{ and } j \in S_{m-1}; \\ 0, & \text{otherwise}. \end{cases}
\]
Given total payments \( P^\text{(m)}_{i} \) and a share price \( \pi^\text{(m)} \), let
\[
A^\text{(m)}_i = c^\text{(m)}_i + \pi^\text{(m)}_i \sum_{k \neq i} P^\text{(m)}_{ik} 
\]  
(28)
denote the value of node \( i \)'s remaining cash and seized and returned collateral. The set of nodes that default in the \( m \)th round is given by
\[
D^\text{(m)} = \left\{ i : i \in S_{m-1}, \text{ and } A^\text{(m)}_i < \sum_{k \neq i} P^\text{(m)}_{ik} \right\}. 
\]  
(29)
Set
\[
a^\text{(m)}_{ij} = \left[ \frac{P^\text{(m)}_{ij} - \pi^\text{(m)}_i \Delta^\text{(m)}_j}{\sum_{k \neq i} P^\text{(m)}_{ik} - \pi^\text{(m)}_i \Delta^\text{(m)}_k} \right].
\]  
Clearing payments in round \( m \) must satisfy
\[
P^\text{(m)}_{ij} = a^\text{(m)}_{ij} \wedge \left[ \frac{\pi^\text{(m)}_i \Delta^\text{(m)}_j + a^\text{(m)}_{ij} A^\text{(m)}_i}{1} \right].
\]  
(30)
The number of shares of returned collateral liquateded by node \( i \) is given by
\[
\Gamma^\text{(m)}_i = \frac{1}{\pi^\text{(m)}} \sum_{j} a^\text{(m)}_{ij} \left( c^\text{(m)}_i + \sum_{k \neq i} P^\text{(m)}_{ik} \right). 
\]  
(31)
With the totals
\[
\Gamma^\text{(m)} = \sum_i \Gamma^\text{(m)}_i \text{ and } \Delta^\text{(m)} = \sum_i \sum_j \Delta^\text{(m)}_{ij} (p^\text{(m)}, \pi^\text{(m)}),
\]
the amount of collateral liquateded in round \( m \) drives the price to
\[
\pi^\text{(m)} = G\left( \pi^\text{(m-1)} + \Gamma^\text{(m)} + \Delta^\text{(m)} \right) = \pi^\text{(m-1)} e^{-\gamma \Gamma^\text{(m)} + \Delta^\text{(m)}}. 
\]  
(32)

**Proposition 4.1.** For any levels of derivatives values \( (\sigma_{ij}, i, j \in N) \), there exist clearing payments and prices \( (p^\text{(m)}, \pi^\text{(m)}), m = 1, \ldots, N \).

With the benefit of this result, we can revisit Example 4.1. Nodes 1 and 3 do indeed default in round 1; their first-round payments are \( P^\text{(1)}_{12} = 8/3 \) and \( P^\text{(1)}_{13} = 4/3 \). The default of node 1 creates a new payment obligation \( \sigma_{12} \) for node 2; the second-round payments are \( P^\text{(2)}_{12} = \sigma_{12}, P^\text{(2)}_{13} = 4/3, \) and \( P^\text{(2)}_{13} = 2/3 \). At the end of the second round, all payment obligations have been met.

### 5. Bankruptcy Stays and Selective Termination

The pros and cons of OTC derivative contract termination at bankruptcy have been debated since the 1990s, and the matter has received renewed attention since the failure of Lehman Brothers. We provide some background before adapting our model to consider some of the key trade-offs.

Most creditors in bankruptcy are subject to a stay that prevents them from seizing assets of a bankrupt entity. This provision is intended to improve the chances that the debtor returns to viability or to maximize the value of the debtor’s assets to repay creditors. Derivatives and certain other financial contracts have long been exempt from these stays. As explained in chapter 9 of Skeel (2010), the exemption was introduced to reduce the risk of spillovers upon the failure of a financial firm by giving special protections to derivatives counterparties.

Since the failure of Lehman Brothers, regulators have come to have a different perspective, seeing termination rights as potentially destabilizing. Fleming and Sarkar (2013) report that Lehman’s derivatives counterparties selectively terminated contracts when they stood to gain but maintained contracts when termination would have resulted in a payment to Lehman, a practice often referred to as cherry-picking. The Federal Deposit Insurance Corporation (2011) reports that contract terminations by Lehman’s counterparties caused market disruptions and left Lehman exposed to greater market risk. In 2017, the Federal Reserve (2017) adopted rules placing some limits on termination of derivatives or, more precisely, qualified financial contracts (QFCs). In explaining the need for an automatic stay on terminations, the Federal Reserve (2017) notes the risk of a “chain reaction” of failures and the risk of fire sales from the liquidation of large volumes of collateral assets.

To capture the phenomenon of creditor cherry-picking, which we refer to as selective termination, we modify the framework of Section 4. We assume that, upon the failure of node \( i \), all amounts \( \sigma_{ij} \) become due, but amounts \( \sigma_{ij} \) do not come due (unless node \( j \) also defaults). In other words, if node \( j \) survives when node \( i \) fails, node \( j \) terminates contracts that trigger obligations from node \( i \) but not contracts that trigger obligations to node \( i \). We compare defaults and payment shortfalls under full termination and selective termination.

### 5.1. Full vs. Selective Contract Termination

We detail the full termination model in Section 4. In contrasting the two scenarios, we use the following notation for the FT and ST models:
\[
P^\text{(f)}_{ij}, P^\text{(s)}_{ij}, \pi^\text{(f)} : \text{payments, payment obligations, and collateral prices in the FT model;}
\]
\[
\sigma^\text{(f)}_{ij}, \sigma^\text{(s)}_{ij}, \tau^\text{(f)} : \text{payments, payment obligations, and collateral prices in the ST model;}
\]
the evolution of the ST model is identical to that of the FT model with one exception: in place of (26), the payment obligations become
\[
\tau^\text{(s)}_{ij} = \begin{cases} 
\sigma^\text{(s)}_{ij} - \pi^\text{(s)}_{ij}, & i \in D^\text{(s-1)} \text{ and } i, j \in S_{m-2}; \\
0, & i \text{ or } j \notin S_{m-2}.
\end{cases}
\]  
(33)
Only the first case has changed: the obligation \( v_j \) is added only if \( i \) defaults.

By construction, the FT and ST models are identical in round 1, so they admit the same sets of clearing payments \( p^{(1)}_{ij} \) and \( q^{(1)}_{ij} \). Moreover, under selective termination, a node that survives round 1 has no payment obligations in round 2, so no defaults occur after round 1, and the process terminates at the end of round 2. By the argument used for Proposition 4.1, we have the following result:

**Proposition 5.1.** For any levels of derivatives values \( \{v_{ij}, i, j \in N\}, \) there exist clearing payments and prices \( (q^{(1)}, \tilde{p}^{(1)}) \) and \( (q^{(2)}, \tilde{p}^{(2)}) \). The feasible clearing payments and collateral prices in round 1 under the ST and FT models coincide, and thus, so do the first-round payment shortfalls. The ST default set is a subset of the FT default set.

Because the two models agree in the first round, they produce the same defaults in round 1, and because the ST model has no subsequent defaults, full termination always produces at least as many defaults as selective termination. The comparison of payment shortfalls is less clear. Full termination creates additional payment obligations and, thus, more opportunities for payments to fall short, but full termination can also increase the flow of payments, potentially offsetting the first effect.

To formulate the comparison precisely, we use the notation in (29) to denote the total default sets for the two models (full and selective termination) by

\[
D^f = \bigcup_{i=1}^{m} D^{(i)}_i, \quad D^s = D^{(1)}.
\]

We define payment shortfalls for the two models as differences between payments due and payments made:

\[
U^f = \sum_{i \in D^f} (v_i + p_i - \tilde{q}_i), \quad U^s = \sum_{i \in D^s} (v_i + p_i - q_i),
\]

where \( p_i = \sum_{j=1}^{w} p^{(1)}_{ij} + q_j = q^{(1)}_i + q^{(2)}_i \), and \( v_i = \sum_{j=1}^{m} v_{ij}. \)

These shortfall measures are calculated relative to payments currently due, which is the relevant focus in a moment of market stress or a crisis following defaults, and ignore future obligations. The total payments due in the FT case are always at least as large as in the ST case, but the comparison of shortfalls can go either way as the following example illustrates.

**Example 5.1.** Consider the linear network of Figure 7 with nodes labeled 0, 1, \ldots, \( N+1 \). Nodes \( i = 1, \ldots, N \) have payments due to their successor nodes, \( \tilde{p}_{i,i+1} = d \). Node 0 has a potential obligation \( v_0 = d \). Node 0 holds \( c_0 \) in cash; no other nodes have cash. In round 1, nodes 1, \ldots, \( N \) default. Under selective termination, nothing more happens; none of the \( Nd \) payments due are made, so \( L^s = Nd \). Under full termination, the default of node 1 triggers termination of the contract between nodes 0 and 1, creating a payment obligation \( p^{(1)}_{01} = v_0 = d \). Node 0 pays \( (d \land c_0) \) to node 1, and this amount is passed through all downstream nodes. The payment shortfall becomes \( U^f = (N+1)d - (N+1)(d \land c_0) = (N+1)(d - c_0) \). By varying the parameters \( N \) and \( d \), we can make \( U^f \) as small as \( L^s \) arbitrarily large, but we can also make \( U^f - L^s \) arbitrarily large.

This example suggests the following properties: if full termination does not increase the set of defaults, then it (weakly) lowers the payment shortfall compared with selective termination. Unless a node is exactly on the boundary of default, a sufficiently small increase in payment obligations will not push it into default. Thus, for sufficiently small \( v_{ij} \), we have \( U^f \leq U^s \), but by increasing some \( v_{ij} \), we can make \( U^f - L^s \) arbitrarily large as long as \( d = 0 \) and \( i \) does not. The key to comparing payment shortfalls is to understand the magnitudes of derivatives liabilities relative to node distances from their default boundaries. We interpret these relative magnitudes as measures of derivatives leverage.

As a first step in formalizing these ideas, we show that full termination increases payments when collateral value is constant. Recall that the ST and FT models coincide in round 1.

**Lemma 5.1.** Suppose the collateral price is constant \( \tilde{p} = \tilde{q} = 1 \). Then, full and selective termination models satisfy

\[
\tilde{p}^{(2)}_{ij} \geq q^{(2)}_{ij}, \quad \text{for all } i, j \in N,
\]

taking the smallest or largest clearing payments under each model.

In the proof of the lemma, we confirm that the smallest and largest clearing payments are well-defined for the two models. We compare these extremal solutions to account for the possibility of nonuniqueness. If each model has unique second-round clearing payments, then (36) holds directly. In light of these considerations, we compare models under payments satisfying either of the following conditions:

\[
p^{(1)}_i = q^{(1)}_i, \quad \tilde{p}^{(1)}_i = \tilde{q}^{(1)}_i, \quad \text{and } p^{(2)}_i \text{ and } q^{(2)}_i \text{ are the largest second round payments;}
\]

\[
p^{(1)}_i = q^{(1)}_i, \quad \tilde{p}^{(1)}_i = \tilde{q}^{(1)}_i, \quad \text{and } p^{(2)}_i \text{ and } q^{(2)}_i \text{ are the smallest second round payments.}
\]

When the ST and FT models produce the same default sets, the additional termination obligations under
the FT model must be fully met; otherwise, they would trigger additional defaults. Combining this observation with Lemma 5.1 yields the following.

**Corollary 5.1.** Suppose either \((37)\) or \((38)\) holds and the collateral price is constant. If \(D^* = D^*\), then \(L^* \leq L^\dagger\).

To build on this observation, we develop the notion of derivatives leverage introduced informally earlier. Let \(e_i\) denote the net worth of node \(i\) at the end of round 1:

\[
e_i = c_i + \pi_i^{(1)} r_i + \sum_{k \neq i} \pi_{ik}^{(1)} - \sum_{k \neq i} \pi_{ki}^{(1)}.
\]

(39)

**Proposition 5.2.** Suppose \((37)\) or \((38)\) holds. Full termination reduces payment shortfalls in the sense that \(L^* \leq L^\dagger\) under the "aggregate derivatives leverage" condition

\[
\sum_{(i,j) \in \Gamma(1)} v_{ij} \leq \sum_{(i,j) \in \Gamma(0)} e_{ij}
\]

(40)

if either (i) the collateral price is constant \(\pi = \pi_i = 1\) or 
(ii) there is no excess collateral, meaning that \(m_0 \leq p_{ij}\), \(\forall i, j \in \mathcal{N}\).

We call (40) a derivatives leverage condition because it compares derivatives liabilities on the left with a measure of equity on the right. A simple sufficient condition for (40) is \(v_{ij} \leq e_i\) for all \(i \notin D(0)\). This condition applies to Example 5.1. Only node 0 survives the first round, and its net worth is \(e_0 = c_0\). If \(c_0 \geq v_0 = d\), then \(L^* = 0\), and thus, \(L^* \leq L^\dagger\). Condition (40) also holds when \(D^* = D^0\)—the case considered in Corollary 5.1.

To add some qualitative context to (40), we note that postcrisis capital and liquidity regulations have significantly increased bank capital and liquidity levels. According to the U.S. Department of the Treasury (2017a), large U.S. banks hold nearly 24% of their assets in high-quality liquid assets, such as cash and U.S. Treasury securities. We know from bank regulatory reporting that derivatives transactions constitute a small portion of the balance sheets of even the largest U.S. bank holding companies. These considerations suggest that the cash demands from contract terminations are unlikely to topple an otherwise solvent bank, favoring full termination over selective termination.

In defining the shortfall measure \(L^\dagger\) in (35), we have not included contracts that have positive value \(v_{ij}\) for a failed node \(i \in D^\dagger\), where \(k \neq D^\dagger\) is a surviving node: under selective termination, these contracts are not terminated, and the payment obligations are not accelerated. To include these quantities as payment shortfalls, we can define

\[
L^{\dagger*} = L^\dagger + \sum_{(i,j) \in \Gamma(0)} v_{ij}.
\]

As \(L^* \leq L^{\dagger*}\), we clearly have \(L^* \leq L^{\dagger*}\) under the conditions of Proposition 5.2.

The following modification of Example 5.1 highlights the effect of collateral fire sales on the total payment shortfalls under full and selective termination.

**Example 5.2.** We modify the linear network of Example 5.1 by setting \(c_0 = 0\) to remove node 0's cash and introducing collateral shares \(m_0 > 0\) posted by node 0 to node 1. With a fixed collateral price \(\pi = 1\) and \(m_0 = c_0\), the shortfalls in this network are identical to those in Example 5.1: following the first-round default of nodes 1, . . . , N, the collateral \(m_0\) is returned to node 0, and node 0 pays \(d \wedge m_0\) to node 1 in the case of full termination; it pays nothing in the case of selective termination. Suppose \(m_0 > d\) so that node 0 could meet its obligation in the FT case if \(\pi = 1\), resulting in \(L^* = 0\). If the collateral is illiquid, its value drops when node 0 sells collateral shares. Applying (31) and (32), we find that the amount liquidated is \(\Gamma(2) = m_0 \wedge d \wedge \pi(2) = \pi(2) - \frac{\pi(2)}{\pi(1)} m_0 \wedge d \wedge \pi(2)\) if \(m_0 < d \wedge \pi(2) m_0\). If \(m_0 \geq d \wedge \pi(2) m_0\), then \(\Gamma(2) = m_0\) (all shares are liquidated), and the amount node 0 pays to node 1 is \(\pi(2) m_0 \wedge d \wedge \pi(2) < d\). The resulting shortfall is \(L^* = (N + 1) d \wedge m_0\)

whereas \(L^* = N d\). We, thus, have \(L^* \geq L^\dagger\) for sufficiently large \(d\). In other words, the illiquidity of the collateral can reverse the order of the two shortfalls.

**5.2. Comparison with No Termination**

If we set aside the option for a bankrupt node to reject certain contracts, we can model an automatic stay by supposing that no payments are accelerated at default. This no-termination scenario is equivalent to setting all \(v_{ij} = 0\), which reduces to the model of Sections 2 and 3. We compare no termination with full and selective termination.

We begin with the comparison between full and no termination. We continue to use \(p_{ij}\) to denote clearing payments under the full termination protocol; in this section, we use \(q_{ij}\) to denote clearing payments with no accelerated payments. The two scenarios coincide in round 1, and without contract terminations, no defaults occur after round 1, and no payments are made after round 2. We compare payments under the assumption that \((37)\) or \((38)\) holds, applying these conditions to the new \(q_{ij}\). The total payment shortfall in the no-termination scenario is given by

\[
L^n = \sum_{(i,j) \in \Gamma(0)} (p_{ij} - q_{ij}),
\]

with \(q_{ij} = q_{ij}^{(1)} + q_{ij}^{(2)}\).

**Proposition 5.3.** Suppose that \((37)\) or \((38)\) holds. Suppose there is no excess collateral, meaning that \(m_{ij} \leq m_{ij}\), \(\forall i, j \in \mathcal{N}\).
Then, full termination results in the default set as no termination and lower payment shortfalls, \( U \leq L^a \), if the following condition holds:

\[
\varphi_i = \begin{cases} 
\varphi_{i}, & i \notin D^{(1)} \\
\sum_{k \in D^{(1)}} \varphi_{k}, & i \in D^{(1)} 
\end{cases}, \quad (41)
\]

The comparison of full and no termination presents trade-offs similar to those in Section 5.1. Contract termination creates additional payment obligations, but it can also increase the flow of cash to meet payment obligations. As in Proposition 5.2, the condition in (41) can be interpreted as a constraint on derivatives leverage.

If we take the view that counterparties of a failed node have to replace their contracts, then we may define

\[
L^a = \sum_{i \in D^{(1)}} (v_i + \bar{p}_i - \bar{q}_i).
\]

This shortfall measure includes the total value \( v_i \) of node \( i \)'s outstanding contracts and, thus, the replacement cost imposed on node \( i \)'s counterparties. Proposition 5.3 clearly applies with \( L^a \) replaced by \( L^a \).

For the comparison of selective termination (creditor cherry-picking) and no termination, we have the following simpler result. Recall that the two models coincide in round 1.

**Proposition 5.4.** Suppose there is no excess collateral, meaning that \( m_0 \leq \bar{p}_i \), \( \forall i \in N \). For any common set of first-round payments under selective termination and no termination, we have \( D^a = D^s \) and \( L^a \leq L^s \).

The network is unlikely to have excess collateral following a large shock. But, if we drop the assumption of no excess collateral, the comparison of shortfalls could go either way. For example, suppose in Figure 7 that node 0 has a payment obligation to some other node \( A \) that it cannot meet, and suppose node 0 has posted excess collateral to node \( A \). That excess collateral is returned to node 0 in round 2. If node 1 selectively terminates its contract with node 0, this creates a new payment obligation \( \bar{m}_0 \), potentially increasing the systemic shortfall. But, if the excess collateral returned to node 0 is large, the contract termination leads to additional payments by nodes 1 through \( N \), potentially reducing the systemicwide shortfall.

**6. Concluding Remarks**

This paper introduces a framework to study contagion in collateralized financial networks and to analyze the effects of the contract termination rules that control access to collateral. We compare alternative scenarios through their impact on the set of nodes that default and the total payment shortfall. In a collateralized network, the failure of one firm may improve the ability of other firms to meet their obligations.

We show that this phenomenon makes the problem of determining clearing payments ill posed. We resolve this difficulty and arrive at a well-defined set of clearing payments by carefully specifying the timing of payments and collateral liquidation.

It is interesting to view our analysis through a regulatory lens. In its comparison of margin requirements and capital requirements, the Basel Committee on Banking Supervision and International Organization of Securities Commissions (2015, p. 4) writes that, “margin can be seen as offering enhanced protection [in comparison with capital] against counterparty credit risk provided that it is effectively implemented. In order for margin to act as an effective risk mitigant, it must be (i) accessible when needed and (ii) provided in a form that can be liquidated rapidly and at a predictable price even in a time of financial stress.” Our analysis of fire sales reinforces the second point, and our baseline model of collateral presupposes that collateral is accessible at default. Our results also show that, even when (i) and (ii) hold, collateral is not guaranteed to improve financial stability. Depending on how collateral is allocated to counterparties, it can increase or decrease defaults and payment shortfalls.

For instance, we show that committing excess collateral may increase risks to financial stability. This result is applicable with derivative contracts in which initial margin can lead to collateral levels in excess of current payment obligations. Moreover, a comparison of alternative policies on collateral seizure requires consideration of a firm’s positions in a network of payment obligations and cannot be made by considering a firm in isolation.

The same point—that the network matters—applies to the comparison of alternative rules on contract termination. The debate over stays on contract termination upon the failure of a firm gained renewed attention after the failure of Lehman Brothers. Policies adopted earlier argued that financial stability required protecting the termination rights of surviving counterparties; more recently, the regulatory consensus argues that financial stability requires limiting these rights. Our analysis considers systemic losses and defaults under alternative assumptions about contract termination.

These comparisons require analyzing the network; our results and examples show that none of the termination scenarios we consider is uniformly better than the others in stemming losses.

We are able to make a stronger statement under a constraint on derivatives leverage in the network. When this condition holds, full termination results in lower payment shortfalls than selective termination or no termination. A reduction in derivatives leverage is consistent with postcrisis increases in bank capital and a general decline in over-the-counter derivatives, adding to the relevance of the condition we introduce.
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Appendix A. Supporting Analysis

A.1. Proof of Lemma 2.1

Write $S_1 = \{ k : p_{k1}^{(1)} < p_{a} \}$ for the set of nodes to which node $i$ did not meet its payment obligations in round 1. If $j \notin S_1$, then $a_{j}^{(2)} = 0$, so we may suppose $j \in S_1$; in particular, $S_1$ is nonempty. The sums in the denominators of $a_{j}^{(1)}$ and $a_{j}^{(2)}$ can be restricted to $k \in S_1$. We prove the result in the more general setting of Section 3, for which $\pi^{(2)} = \pi^{(3)} = 1$ is a special case. For all $k \in S_1$,

$$p_{k}^{(1)} = \pi^{(1)} m_{k} + a_{k}^{(1)} \Lambda_{k}^{(1)}.$$

Making this substitution in (11) and (13) and letting $K_i$ denote the denominator of $a_{j}^{(1)}$ yields

$$a_{j}^{(2)} = \frac{\pi^{(1)} m_{j} - a_{j}^{(1)} \Lambda_{j}^{(1)}}{\sum_{k \in S_1} \pi^{(1)} m_{k} - a_{k}^{(1)} \Lambda_{k}^{(1)}} \frac{\sum_{k \in S_1} p_{k}^{(1)} - \pi^{(1)} m_{k} - a_{k}^{(1)} \Lambda_{k}^{(1)}}{\sum_{k \in S_1} (p_{k}^{(1)} - \pi^{(1)} m_{k} - a_{k}^{(1)} \Lambda_{k}^{(1)})}.$$

$$= \frac{\sum_{k \in S_1} p_{k}^{(1)} - \pi^{(1)} m_{k} - a_{k}^{(1)} \Lambda_{k}^{(1)}}{\sum_{k \in S_1} (1 - \Lambda_{k}^{(1)})}.$$

$$= a_{j}^{(1)}.$$

A.2. Proof of Proposition 2.1

The existence of clearing payments in this setting is a special case of the more general claim in Proposition 3.1, which we prove separately. Here we prove (15). The second-round payments in (14) satisfy $p_{k}^{(2)} = p_{k}^{(1)} - p_{k}^{(1)}$, using (11), so $p_{k}^{(1)} + p_{k}^{(1)} - p_{k}^{(1)}$. Summing over $j$, we get

$$p_{k}^{(1)} + p_{k}^{(2)} = \bar{p}_{k}. \quad \text{(A.1)}$$

We next show that the second term in (15) is also an upper bound:

$$p_{k}^{(1)} + p_{k}^{(2)} \leq \bar{p}_{k} + \sum_{k \in S_1} p_{k}^{(1)} + \sum_{k \in S_1} m_{k}. \quad \text{(A.2)}$$

The second-round payments in (14) have exactly the structure of Eisenberg and Noe (2001) clearing payments (4) with $r_i$ in (14) playing the role of $c_i$ in (4) and payment obligations given by $p_{k}^{(1)}$. Any node that did not default in the first round has no payment obligations in the second round. As in Eisenberg and Noe (2001), we may write the node totals as

$$p_{k}^{(2)} = p_{k}^{(2)} \wedge \left( r_{i} + \sum_{k \in S_1} p_{k}^{(2)} \right).$$

In other words, the minimum in (14) is either attained by the first term for all $j$ or the second term for all $j$. For $i \in D$, we can rewrite $r_i$ using (5) and (12) as $r_i = \sum_{k \in S_1} (m_{k} - \bar{p}_{k})$. Making this substitution and using (11), we get

$$p_{k}^{(2)} = \left[ \bar{p}_{k} - p_{k}^{(1)} \right] \wedge \left( \sum_{k \in S_1} (m_{k} - \bar{p}_{k}) \wedge \sum_{k \in S_1} p_{k}^{(1)} \right). \quad \text{(A.3)}$$

From (10), we have

$$p_{k}^{(1)} = (m_{k} \wedge \bar{p}_{k}) + a_{k}^{(1)} \Lambda_{k}^{(1)} \wedge \bar{p}_{k}.$$

Summing both sides of (A.4) over $k$ and recalling the definition of $A_{k}^{(1)}$ in (7), we get

$$p_{k}^{(1)} \leq \bar{c}_{i} + \sum_{k \in S_1} m_{k} \wedge \bar{p}_{k}. \quad \text{(A.5)}$$

Adding (A.5) to (A.3) yields (A.2).

In light of (A.1) and (A.2), to prove (15), we need to show that

$$p_{k}^{(1)} + p_{k}^{(2)} < \bar{p}_{k} \Rightarrow p_{k}^{(1)} + p_{k}^{(2)} = \bar{c}_{i} + \sum_{k \in S_1} m_{k} \wedge \bar{p}_{k}. \quad \text{(A.6)}$$

We claim that, if $p_{k}^{(1)} + p_{k}^{(2)} < \bar{p}_{k}$, then equality holds in (A.4), and (10) can be written as

$$p_{k}^{(1)} = (m_{k} \wedge \bar{p}_{k}) + a_{k}^{(1)} \Lambda_{k}^{(1)}.$$

To show this equivalence, we need to consider three cases. (i) If $\bar{p}_{k} \leq m_{k}$, then (6) yields $a_{j}^{(1)} = 0$, so (10) and (A.7) both give $p_{k}^{(1)} = \bar{p}_{k}$, (ii) If $\bar{p}_{k} \geq m_{k} + a_{k}^{(1)} \Lambda_{k}^{(1)}$, then (10) and (A.7) both give $p_{k}^{(1)} = m_{k} + a_{k}^{(1)} \Lambda_{k}^{(1)}$. (iii) The remaining case is $m_{k} < \bar{p}_{k} < m_{k} + a_{k}^{(1)} \Lambda_{k}^{(1)}$, which is equivalent to $0 < \bar{p}_{k} - m_{k} < a_{k}^{(1)} \Lambda_{k}^{(1)}$. In light of the definition of $a_{k}^{(1)}$ in (6), this implies $\sum_{j \in S_1} p_{j}^{(1)} + m_{k} < A_{k}^{(1)}$. But this inequality says that, following the seizure of collateral, node $i$ has sufficient remaining assets to meet all its residual claims, making $p_{k}^{(1)} = \bar{p}_{k}$ for all $k$. Summing over $k$ yields $p_{k}^{(1)} = \bar{p}_{k}$. Thus, under the condition $p_{k}^{(1)} + p_{k}^{(2)} < \bar{p}_{k}$, case (iii) is precluded and (A.7) holds.

Summing over $k$ in (A.7) we get equality in (A.5). Under the condition $p_{k}^{(1)} + p_{k}^{(2)} < \bar{p}_{k}$ in (A.6), the minimum in (A.3) is attained by the second term. Adding this term to the right side of (A.5) yields the claimed result in (A.6).

A.3. Analysis of Free-Termination Model

We begin with a precise formulation of the model of Section 2.5. To capture the nodes’ access to collateral, we need to et

$$A_{i} = c_{i} + \sum_{k \in S_1} (m_{k} \wedge \bar{p}_{k}) \wedge \sum_{k \in S_1} m_{k}. \quad \text{(A.8)}$$

The term $\sum_{k \in S_1} m_{k} \wedge \bar{p}_{k}$ reflects the excess collateral node $i$ can call back from $k$ to make payments to other nodes. Node $i$
defaults if \( A_i < \sum_k \{ p_k - m_k \}^+ \); in other words, default is the failure to meet uncollateralized obligations. We can write the default set as

\[
D^p = \left\{ i : c_i + \sum_k p_k + \sum_k m_k < \sum_k p_k \right\}.
\]  
(A.9)

Upon \( f \)'s default, node \( f \)'s share of any remaining assets is proportional to its residual claim, so

\[
a_q = -\frac{\{ p_f - m_f \}^+}{\sum_k \{ p_k - m_k \}^+}.
\]  
(A.10)

Clearing payments are required to satisfy

\[
p_q = p_q \land \{ m_q + a_q A_q \}.
\]  
(A.11)

Existence of (largest and smallest) clearing payments follows from Tarski's (1955) fixed-point theorem.

To formulate the equivalent Eisenberg–Noe model, define reduced obligations

\[
q_q = \frac{\{ p_q - m_q \}}{\sum_k \{ p_k - m_k \}^+}.
\]  
(A.12)

and increased cash

\[
c_q^+ = c_i + \sum_k \{ m_k - p_k \}^+ + \sum_k \{ m_k \land p_k \}
\]  
(A.13)

The additional cash reflects collateral \( \{ m_k - p_k \}^+ \) recovered by \( i \) and any paying down of obligations to \( i \) using collateral posted by \( k \), \( m_k \land p_k \). Set

\[
A_q^+ = c_i + \sum_k q_k
\]  
(A.14)

and notice that \( a_q \) in (A.10) equals \( \frac{q_q}{\sum_k q_k} \). With no collateral, the standard Eisenberg–Noe condition for clearing payments becomes

\[
q_q = \frac{q_q}{\sum_k q_k} A_q^+.
\]  
(A.15)

We may rephrase the first statement of Lemma 2.2 as saying that payments \( q_q \) satisfy (A.13) if and only if payments \( p_q = q_q \land \{ m_q + a_q A_q^+ \} \) satisfy (A.11).

Proof of Proposition 2.3. For the model with free termination, we can use \( c_q^+ \) in (A.8) to rewrite the clearing condition (A.11) as

\[
p_q = p_q \land \{ m_q + a_q \left( c_i + \sum_k p_k \right) \}.
\]  
(A.16)

This equation has exactly the same form as the first-round clearing payments (10) but with \( c_i \) in (7) replaced by \( c_q^+ \). As \( c_q^+ \geq c_i \), it follows from theorem 3 of Milgram and Roberts (1994) that the largest and smallest fixed points of (A.13) are no smaller than, respectively, the largest and smallest fixed points of (10). In other words, payments with free termination exceed first-round payments in the original model. By comparing (8) and (A.9), we see that \( p_q \geq p_q^{(0)} \) implies \( D^p \subseteq D \): free termination results in fewer defaults.

To compare payment shortfalls in the two models, we use Equation (A.16), proved as follows, and claim that we can replace \( c_i + r_i \) with \( c_q^+ \) to write

\[
\frac{p_q^{(1)}}{p_q^{(0)}} + \frac{p_q^{(2)}}{p_q^{(0)}} = p_q \land \left( \{ m_q + a_q \left( c_i + \sum_k p_k \right) \} \right).
\]  
(A.17)

If \( i \notin D \), then \( p_q^{(0)} = p_q, p_q^{(0)} = 0 \), and there is nothing to show. If \( i \in D \), then the returned collateral is \( r_i = \sum_k \{ m_k - p_k \}^+ \), and indeed \( c_i + r_i = c_q^+ \). Comparison with (A.15) now shows that total payments in the two systems coincide.

Proof of Proposition 2.4. We first derive an expression for total payments \( p_q^{(1)} + p_q^{(2)} \) that is of independent interest. Using Lemma 2.1, we can replace \( a_q \) in (14) with \( a_q \) because, if \( a_q = 0 \), then \( p_q^{(1)} = p_q \); so \( p_q^{(0)} = 0 \) and is unchanged by the replacement. With this substitution and adding (10) and (14), we get

\[
\frac{p_q^{(1)}}{p_q^{(0)}} + \frac{p_q^{(2)}}{p_q^{(0)}} = p_q \land \left( \{ m_q + a_q \left( c_i + \sum_k p_k \right) \} \right)
\]  
(A.18)

This expression does not quite reduce the two rounds to a single round because the returned collateral \( r_i \) is determined after the first round and is not an exogenous parameter. We now turn to the two claims in the proposition.
i. Removing excess collateral has no effect on round 1 payments and no effect on which nodes default. All excess collateral becomes returned collateral in round 2. Replacing excess collateral with cash increases \( \zeta_i \) by exactly the amount it decreases returned collateral \( r_i \), so we can see from (A.16) that pooling excess collateral does not affect the set of clearing payments and, therefore, does not affect payment shortfalls. The increase in \( \zeta_i \) can, however, reduce the default set, which is determined in round 1.

ii. Under proportional collateral, \( \bar{p}_y = m_y \) for \( y = 1, \ldots, n \), so \( a_{ij} = p_i / m_j \), from which it follows that \( m_y = a_{ij} \) with \( m_y = \sum_i m_y \). Moreover, with \( k < 0 \), there is no excess collateral, so no defaulting node receives any returned collateral, meaning that \( a_{ij} \rho_i r_i = 0 \). We may, therefore, write (A.16) as

\[
p_{ij}^{(1)} + p_{ij}^{(2)} = \bar{p}_y \wedge \mathcal{A}_j \left( m_i + \zeta_i + \sum_k \left( a_{ij}^{(1)} + a_{kj}^{(2)} \right) \right).
\]

As the total payments depend on \( m_i \) and \( \zeta_i \) only through their sum, pooling while preserving proportional collateral (increasing \( \zeta_i \) by decreasing \( k < 0 \)) has no effect on payment shortfalls, but as before, increasing \( \zeta_i \) can reduce defaults.

### A.5. Proof of Proposition 3.1

We begin with the analysis of first-round payments. Through an arbitrary ordering of pairs of nodes, we can record the set of payments \( p_{ij} \) in a vector \( \mathcal{P} \), interpret the vector \( \mathcal{P} \) accordingly. If we take any \( \ell \in [0, 1] \) and \( p_{ij} \in [0, \bar{p}_y] \) and plug these variables into the right side of (17)–(22), then the variables on the left side of (17) and (22) return new values of \( \ell \in [0, 1] \) and \( p_{ij} \in [0, \bar{p}_y] \). In other words, expressions (17)–(22) define a mapping \( F: (\pi, \rho^{(3)}) \to (\pi, \rho^{(3)}) \) of \( [0, 1] \times [0, \bar{p}_y] \) to itself.

**Lemma A.1.** \( F \) is monotone increasing.

**Proof.** In (22), we see that \( p_{ij} \) is monotone increasing in \( \mathcal{A}_j \) and, therefore, monotone increasing in \( p_{ij}^{(1)} \), \( k \neq l \). Now, consider \( p_{ij}^{(1)} \) as a function of \( \pi \). If \( i \notin D \), then \( p_{ij}^{(1)} = \bar{p}_y \), and changing \( \pi \) has no effect on \( p_{ij}^{(1)} \). For \( i \in D \), we consider three cases.

1. **Case 1:** Suppose that, \( \mathcal{A}_j < \sum_\ell (\bar{p}_y - \rho^{(3)}) \zeta_i \) and \( \bar{p}_y > \rho^{(3)} \). In this case, (22) yields a right derivative of

\[
\frac{\partial p_{ij}^{(1)}}{\partial \pi} = m_i - \frac{m_y \mathcal{A}_j}{\sum_\ell (\bar{p}_y - \rho^{(3)}) \zeta_i} + \frac{\bar{p}_y - \rho^{(3)}}{\sum_\ell (\bar{p}_y - \rho^{(3)}) \zeta_i} - \frac{m_y \mathcal{A}_j}{\sum_\ell (\bar{p}_y - \rho^{(3)}) \zeta_i}.
\]

The second term on the right is less than \( m_y \) because we have assumed \( \mathcal{A}_j < \sum_\ell (\bar{p}_y - \rho^{(3)}) \zeta_i \). The third term is nonnegative, and the derivative is then as well. A small increase in \( \pi \) yields an increase in \( p_{ij}^{(1)} \).

2. **Case 2:** Suppose that, \( \mathcal{A}_j > \sum_\ell (\bar{p}_y - \rho^{(3)}) \zeta_i \) and \( \bar{p}_y > \rho^{(3)} \). These conditions imply \( p_{ij}^{(1)} = \bar{p}_y \), and they are preserved under a small increase in \( \pi \), so \( \partial p_{ij}^{(1)}/\partial \pi = 0 \).

3. **Case 3:** Suppose that, \( \bar{p}_y \leq \rho^{(3)} \). In this case, \( p_{ij}^{(1)} = \bar{p}_y \), so \( \partial p_{ij}^{(1)}/\partial \pi = 0 \).

As \( \pi \) increases, we may transition into cases 2 or 3. Either transition yields \( p_{ij}^{(1)} = \bar{p}_y \), so monotonicity holds.

It remains to show that (17) makes \( \pi \) on the left an increasing function of all \( p_{ij}^{(1)} \) and \( \pi \) on the right. Monotonicity in \( \pi \) is immediate from the monotonicity of \( G \). An increase in \( p_{ij}^{(1)} \) leads the default set \( D \) to contract or remain unchanged, resulting in a decrease in \( \Delta \) and an increase in \( \pi \).

By Tarski’s (1955) fixed-point theorem, Lemma A.1 implies that \( F \) has a fixed point in \( [0, 1] \times [0, \bar{p}_y] \), which delivers the required clearing payments and collateral price. The following lemma ensures that a fixed point is reached through iterative application of \( F \), starting from the upper boundary of this domain.

**Lemma A.2.** \( F \) is continuous from the right; that is, for any decreasing convergent sequence \( (\pi_n, p_n^{(1)}) \to (\pi^*, p^{(1)*}) \), we have \( F(\pi_n, p_n^{(1)}) \to F(\pi^*, p^{(1)*}) \).

**Proof.** Equations (17)–(22) imply that \( p_{ij}^{(1)} \) is continuous in \( p_{ij}^{(3)} \) and \( \pi \), and (17) implies \( \pi \) on the left is continuous in \( \pi \) on the right. A change in \( p_{ij}^{(1)} \) may produce a change in \( D \) and, thus, a discontinuity in \( \pi \). However, a small increase in \( p_{ij}^{(1)} \) preserves the inequalities defining the default set in (18), leaving \( D \) unchanged and implying right-continuity of \( \pi \) in \( p_{ij}^{(1)} \).

We can now conclude the proof of Proposition 3.1. By Lemma A.1, the iterates of \( F \) starting from \( (1, \bar{p}_y) \) form a decreasing sequence \( F(\pi^{(i)}, p^{(i)}) = (\pi^{(i+1)}, p^{(i+1)}) \). This sequence is bounded below by \( (0, 0) \) and, therefore, has a limit \( (\pi^*, p^{(1)*}) \). Thus, \( F(\pi^*, p^{(1)*}) = (\pi^*, p^{(1)*}) \). However, Lemma A.2 implies that \( F(\pi^*, p^{(1)*}) \to F(\pi^*, p^{(1)*}) \). We conclude that \( F(\pi^*, p^{(1)*}) = (\pi^*, p^{(1)*}) \).

This fixed point provides a first-round price and clearing vector \( (\pi^{(1)}, p^{(1)*}) \). The existence of \( (\pi^{(2)}, p^{(2)*}) \) now follows by a similar but simpler argument. Expressions (23)–(25) define a mapping from \( (\pi^{(2)}, p^{(2)*}) \) to the right \( (\pi^{(2)}, p^{(2)*}) \) on the left. The mapping is clearly monotone increasing and continuous, and it maps \([0, 1] \times [0, \bar{p}_y] \) into itself. It, therefore, has a fixed point, and the fixed point delivers the required solution \( (\pi^{(2)}, p^{(2)*}) \).

### A.6. Proof of Proposition 3.2

The existence of a largest (and smallest) first-round solution follows from Tarski’s (1955) fixed-point theorem and the monotonicity of \( F \) in \( \pi^{(2)} \). By theorem 3 of Milgrom and Roberts (1994), the monotonicity of \( F \) in \( \pi^{(2)} \) implies that the largest and smallest \( \pi^{(2)} \) are smaller than liquid collateral \( (\pi^{(2)} \leq 1) \) than with illiquid collateral \( (\pi^{(2)} \leq 1) \). This implies that the default set is larger with illiquid collateral.

With illiquid collateral, the argument leading to (A.16) yields

\[
p_{ij}^{(1)} + p_{ij}^{(2)} = \bar{p}_y \wedge \left( \pi^{(2)} m_i + a_{ij} \left( \pi^{(2)} + \sum_k \left( p_{ik}^{(1)} + p_{kj}^{(2)} \right) \right) \right),
\]

\[
= \bar{p}_y \wedge \left( \pi^{(2)} m_i + a_{ij} \Lambda_i \right).
\]

With this representation, we may write \( p_{ij}^{(2)} = p_{ij}^{(2)} + p_{ij}^{(1)} + p_{ij}^{(2)} = (F(p_{ij}^{(1)}, p_{ij}^{(2)}), F(p_{ij}^{(1)}, p_{ij}^{(2)})), F(p_{ij}^{(1)}, p_{ij}^{(2)})) \) with \( F \) as in Lemma A.1 and \( F \) defined by (A.19) and (24).

In (A.19), \( \Lambda_i \) is increasing in \( p_{ij}^{(1)} \) and \( p_{ij}^{(2)} \), in \( \pi^{(2)} \), and in \( n_i \), which is increasing in \( \pi^{(2)} \). Moreover, (A.19) has the same form as (22), so the argument in Lemma A.1 shows that \( F \) is monotone increasing. It follows from Tarski’s (1955)
fixed-point theorem that the mapping defined by \( F, \hat{F} \) has a largest and smallest fixed point. By theorem 3 of Milgrom and Roberts (1994), with \( \pi^{(i)}, \pi^{(2)} \leq 1 \) (illiqid collateral), the largest and smallest fixed points are smaller than the largest and smallest fixed points when we set \( \pi^{(1)} = \pi^{(2)} = 1 \) (the case of liquid collateral). As illiquid collateral yields smaller total payments \( p^{(i)} + p^{(0)} \), it yields larger payment shortfalls.

### A.7. Analysis of Collateral-First Protocol

In several places in this paper, we compare our baseline model of Sections 2.2, 2.3, 3.1, and 3.2 with alternative models. In making these comparisons, we use the notation \( q_j \) and \( h \) (with superscripts to distinguish rounds) to denote payments and prices in the alternative model. The specific meaning of these variables is different in different sections as we use this notation to compare our baseline model against different alternatives.

In this section, we use \( q_j^{(1)}, q_j^{(2)} \) and \( h^{(1)}, h^{(2)} \) to denote two rounds of payments and collateral prices under a protocol in which payments precede collateral seizure in round 1. First-round payments are characterized by

\[
q_j^{(1)} = \left\{ \begin{array}{ll}
\frac{p_h}{\pi_j} + a_j^{(1)} A_j^{(1)} + \frac{[p_h - a_j^{(1)} A_j^{(1)}]}{\pi_j} \land h_m & \text{if } i \in D^1; \\
\frac{p_h}{\pi_j} & \text{if } i \notin D^1.
\end{array} \right.
\]  

(A.20)

The first case can be read as follows: node \( i \) makes a partial payment to node \( j \) of \( a_j^{(1)} A_j^{(1)} \); any residual obligation \( \frac{p_h - a_j^{(1)} A_j^{(1)}}{\pi_j} \) is paid from collateral, up to the amount available \( h_m \); the total payment cannot exceed the amount due \( p_h \). The assets \( A_j^{(1)} \) have the same form as in (7) but now with incoming payments \( q_j^{(1)} \); the default set is determined exactly as in (8), but we have labeled it \( D^1 \) to indicate its dependence on the payments \( q_j^{(1)} \). The proportions \( a_j^{(1)} \) are as given in (19) and, thus, reflect the collateral posted. Upon node \( i \)'s default, the shares of collateral seized and sold by node \( j \) become

\[
A_j^{(0)} = \left\{ \begin{array}{ll}
\frac{p_h}{\pi_j} & \text{if } i \in D^1; \\
0 & \text{if } i \notin D^1.
\end{array} \right.
\]  

(A.21)

The first case in (A.21) captures the feature that a failed node's collateral can be seized and liquidated only after its other assets are exhausted. Set \( \Delta_j = \sum_j A_j^{(0)} \). As before, the price impact function \( h^{(1)} = e^{\cdot m_{\pi}} \) determines the equilibrium asset price. Once first-round payments and \( h^{(1)} \) are determined, second-round payments \( q_j^{(2)} \) and \( h^{(2)} \) are characterized by (23)-(25), just as before. The proof of Proposition 3.1 can be used to show the existence of first- and second-round clearing payments and prices \( q_j^{(1)}, h^{(1)} \) and \( q_j^{(2)}, h^{(2)} \).

In the case of liquid collateral, \( \pi = 1 \), total payments in the original (collateral-first) model and the alternative (payments-first) model are the same:

### Proposition A.1

Assume that collateral is posted in cash so that there are no fire-sale effects. If \( p_j^{(1)} + r_j^{(2)}, j \in N_j \), are total clearing payments for the original model, then \( q_j^{(1)} = p_j^{(1)} \) and \( q_j^{(2)} = p_j^{(2)} \), \( j \in N_j \), are total clearing payments in the alternative model with delayed collateral seizure. Thus, with liquid collateral, the two protocols yield the same default set and the same payment shortfalls.

### Proof of Proposition A.1

We first show that, if all incoming payments to node \( i \) agree under the two models, \( q_j^{(1)} = p_j^{(1)} \), \( k \in N_i \), then outgoing payments \( p_j^{(2)} \) given by (10) and \( q_j^{(2)} \) given by (A.20) agree for all \( j \in N_j \).

If all incoming payments to node \( i \) agree under the two models, then the two models yield the same \( A_i^{(1)} \) and \( i \notin D^1 \) if and only if \( i \notin D^1 \). If \( i \notin D^1 \), then \( q_j^{(1)} = \frac{p_h}{\pi_j} + p_j^{(0)} \) for all \( j \).

Suppose \( i \notin D^1 \). If \( p_j^{(1)} \leq A_j^{(1)} \), then (10) and (A.20) both evaluate to \( p_h \). If \( p_j^{(1)} > A_j^{(1)} \), then (A.20) evaluates to

\[
q_j^{(1)} = \frac{p_h}{\pi_j} \land \left[ \left( q_j^{(1)} - a_j^{(1)} A_j^{(1)} \right) \land h_m \right],
\]

which agrees with (10). Thus, \( q_j^{(1)} = p_j^{(1)} \) for all \( j \).

We now turn to the second-round payments. We assume that all first-round payments agree under the two models, and we assume that all second-round payment obligations agree, \( q_j^{(2)} = p_j^{(2)} \) for all \( k \in N \), and we show that this implies that \( q_j^{(2)} = p_j^{(2)} \) for all \( j \in N_j \).

If \( i \notin D^1 \), then node \( i \) has no second-round payment obligations, so \( p_j^{(2)} = p_j^{(0)} = 0 \) for all \( j \). Suppose \( i \in D^1 \). If the amount of returned collateral \( r_i \) is the same in the two models, then the payments made by node \( i \) under the two models agree because they are determined by (14).

The only remaining case to consider is the possibility that the two models may produce different quantities of returned collateral \( r_i \). We see from (12) that the quantities of liquidated collateral \( \Delta_i \) in (8) and \( \Delta_j^{(1)} \) in (A.21) must then differ for some \( j \).

We always have \( \Delta_j^{(0)} < \Delta_j^{(1)} \leq \Delta_j \), so for \( i \), and (A.20) differs, we must have \( \Delta_j^{(0)} < \Delta_j^{(1)} < \Delta_j \), and thus, \( \Delta_j^{(0)} < \Delta_j^{(1)} \).

However, this inequality states node \( i \) has enough assets to meet all its round 1 obligations under the collateral-first protocol, so \( p_j^{(0)} = p_j^{(1)} \); hence, \( q_j^{(1)} = p_j^{(0)} \), implying that \( p_j^{(2)} = p_j^{(1)} = 0 \).

### A.8. Proof of Proposition 4.1

The first round with contract termination is identical to the first round in Section 3.1, so the existence of \( p_i^{(1)} \) follows from Proposition 3.1. For \( m \geq 2 \), we claim that \( \Delta_j^{(0)} = 0 \). In light of (22), it suffices to consider the case \( i \in D_j^{(0)} \), \( j \in S_m-1 \). But, if \( i \in D_j^{(0)} \), then \( i \) survived the first \( m-1 \) rounds, so \( i \in S_{m-1} \). In this case, (26) yields \( p_i^{(0)} = 0 \), and (27) yields \( \Delta_j^{(0)} = 0 \). With all \( \Delta_j^{(0)} = 0 \), the mapping defined by Equations (26)-(32) becomes a special case of the mapping in Lemma A.1 with one modification, which is the inclusion of the returned collateral value \( \pi^{(m)} \) in \( \Delta_i^{(m)} \) in (28).

The monotonicity proved in Lemma A.1 holds with this
modification: the only step affected is (A.17), in which we
pick up an additional positive term from the monotonicity
of \( A_{ij}^{(m)} \) in \( \pi^{(m)} \). The existence of \((p^{(m)}, \pi^{(m)})\) follows as in
Proposition 3.1.

A.9. Proof of Proposition 5.2

Proof of Lemma 5.1. The second-round payment obligations
in the FT model are given by

\[
p_{ij}^{(2)} = \begin{cases} 
q_{ij} + p_j - p_i, & \text{i} \neq j, i \in D^{(1)}; \\
0, & \text{otherwise}.
\end{cases}
\]  

(A.22)

This follows from (26) with \( m = 2 \), recalling that all nodes
are in \( S_0 \). Under the ST model, \( q_{ij}^{(1)} = p_{ij}^{(1)} \), and the second-
round payment obligations are given by

\[
q_{ij}^{(2)} = \begin{cases} 
q_{ij} + p_j - p_i, & \text{i} \neq j, i \in D^{(1)}; \\
0, & \text{otherwise}.
\end{cases}
\]  

(A.23)

because the amount \( q_{ij} \) becomes due only if \( i \) defaults. Com-
paring (A.22) and (A.23), we can write \( p_{ij}^{(2)} \geq q_{ij}^{(2)} \), \( \forall i, j \in N \).

Second-round payments under full termination are given by

\( (30) \) with \( m = 2 \). In the proof of Proposition 4.3, we show that \( \Delta^{(m)} = 0 \) for \( m \geq 2 \), so (30) yields

\[
p_{ij}^{(2)} = p_{ij}^{(2)} \land q_{ij}^{(2)} + \sum_{k \in \mathcal{K}_i} p_{ij}^{(2)} 
\]  

(A.24)

taking \( d_{ij}^{(2)} = 0 \) if node \( i \) has no second-round obligations.

Under selective termination,

\[
q_{ij}^{(2)} = q_{ij}^{(2)} \land d_{ij}^{(2)} = \left( d_{ij}^{(2)} + q_{ij}^{(2)} \right) / \sum_{k \in \mathcal{K}_i} d_{ij}^{(2)}.
\]  

(A.25)

The cash amounts \( c_{ij}^{(2)} \) and returned collateral \( r_{ij}^{(2)} \) are indeed
equal in these two expressions because the two models
agree in round 1. Equations (A.24) and (A.25) represent \( p_{ij}^{(2)} \) and \( q_{ij}^{(2)} \) as fixed points of a common mapping, parameter-
ized by \( \bar{p}_{ij}^{(2)} \) and \( \bar{q}_{ij}^{(2)} \) in the first case and by \( \bar{q}_{ij}^{(2)} \) and \( \bar{d}_{ij}^{(2)} \) in
the second case. Moreover, \( p_{ij}^{(2)} \) in (A.24) is an increasing function of \( \bar{p}_{ij}^{(2)} \) and \( \bar{q}_{ij}^{(2)} \), and \( \bar{q}_{ij}^{(2)} \) in (A.25) is an increasing function of \( \bar{q}_{ij}^{(2)} \) and \( \bar{d}_{ij}^{(2)} \).

We have shown that \( p_{ij}^{(2)} \geq q_{ij}^{(2)} \); we now claim that \( d_{ij}^{(2)} \geq q_{ij}^{(2)} \) for all \( i, j \in N \). If \( i \neq N \), then we see from (A.22) and
(A.23) that \( p_{ij}^{(2)} \geq q_{ij}^{(2)} \) for all \( k \), so \( q_{ij}^{(2)} = d_{ij}^{(2)} \), and if \( i \neq D^{(1)} \), then \( d_{ij}^{(2)} = 0 \). The ordering in (36) now follows from theorem 3 of
Milgrom and Roberts (1994) for the smallest and largest
fixed points of the two models, which also ensures the existence of these extremal fixed points.

Proof of Proposition 5.2. From (34), we have \( D' \subset D' \), and as
noted there, \( D' = D' \). To compare payment shortfalls, we use
(35) to write

\[
L' - U' = \sum_{i \in D' \setminus D'} \left( p_i - q_i \right) - \sum_{i \in D' \setminus D'} \left( v_i + p_i - p_i \right) 
\]  

(A.26)

\[
= \sum_{i \in D' \setminus D'} \left( \sum_{j \in \mathcal{N}} p_{ij}^{(1)} - q_{ij}^{(1)} \right) - \sum_{i \in D' \setminus D'} \left( v_i + p_i - p_i \right) 
\]  

(A.27)

Equation (A.26) uses two properties: the first-round pay-
ments agree, \( p_{ij}^{(1)} = q_{ij}^{(1)} \), for all \( i \neq i \) and under selective termi-
nation, there are no payments after the second round, so \( q_{ij}^{(1)} = 0 \) for \( \ell > 2 \). Equation (A.27) follows because, if \( i \neq D' \), then
node \( i \) must have met its first-round payment obligations,
so \( p_i = p_{ij}^{(1)} \).

Under full termination, if node \( i \) survives round 1 but defaults in a subsequent round (i.e., \( i \neq D' \)), then its payments \( \sum_{j \in \mathcal{N}} p_{ij}^{(1)} \) must be at least as large as its net worth \( v_i \) defined in (39); if a node paid out less than its net worth, it
would not default. From (A.27), we get

\[
L' - U' \geq \sum_{i \in D' \setminus D'} \left( \sum_{j \in \mathcal{N}} p_{ij}^{(1)} - q_{ij}^{(1)} \right) 
\]  

(A.28)

\[- \sum_{i \in D' \setminus D'} \left( v_i - c_i \right).\]

In case (i), Lemma 5.1 applies, so \( p_{ij}^{(2)} \geq q_{ij}^{(2)} \), and the first
term on the right in (A.28) is positive. Under condition (40),
we conclude from (A.28) that \( L' \geq U' \). For case (ii), we note
that, if \( m_{j} \leq p_j \) for all \( j \in \mathcal{N} \), then, under selective termi-
nation, nodes that default in round 1 do not have any col-
lateral returned and do not receive any subsequent payments,
so they cannot make any subsequent payments; thus, \( q_{ij}^{(1)} = 0 \)
in (A.26), so (40) again implies \( L' \geq U' \).

Proof of Proposition 5.3. Under full termination and the
condition \( v_i \leq c_i \), \( \forall i \notin D' \), a node that survives the first round can meet its second-round obligations because these obligations \( v_i \) do not exceed the node’s net worth \( c_i \). Thus, \( D' = D' \), and without
further defaults, \( p_{ij}^{(1)} = 0 \) for all \( \ell > 2 \).

In the no-termination scenario, \( D' \) is also the default set
because the two models agree in round 1. In the absence of
excess collateral, a node that defaults in round 1 has no collateral
returned. If no contracts are terminated, then such a
node has no inflow of cash and, therefore, cannot make any
further payments: \( q_{ij}^{(1)} = 0 \) for \( \ell \geq 2 \). If \( i \notin D' \), then
node \( i \) has no payment obligations after round 1, so again \( q_{ij}^{(1)} = 0 \) for \( \ell \geq 2 \). The difference in payment shortfalls is given by

\[
L^* - L' = \sum_{i \in D^* \setminus D'} \left( p_i - q_i \right) - \sum_{i \in D^* \setminus D'} \left( v_i + p_i - p_i \right) 
\]  

(A.29)
The internationally agreed upon principles governing bilateral margin (Basel Committee on Banking Supervision and International Organization of Securities Commissions 2015) specify that it should be “immediately available to the collecting party in the event of the counterparty’s default” (p. 23). They also state that “collateral collected as initial margin from the customer is treated as a customer asset” (p. 24) and “the collected margin must be subject to arrangements that fully protect the posting party to the extent possible under applicable law in the event that the collecting party enters bankruptcy” (p. 8). Our model is designed to capture these principles in simplified form: collateral belongs to the posting party, it can be seized quickly by the collecting party if the posting party defaults, and the return of collateral to the posting party may be delayed by bankruptcy proceedings if the collecting party fails.

In the centrally cleared market, CCPs collect but do not post collateral. Here, too, collateral arrangements ensure immediate access by the CCP in case of a counterparty’s default. Trades with a CCP cannot be unilaterally terminated because the CCP needs to maintain a “matched book” with an offsetting contract for every trade.

### Endnotes

1. Researchers also debate the implications for economic growth of the demand for safe assets as collateral; see Duffie et al. (2015) and Sidanis and Zikes (2012) for estimates of the demand. Our analysis addresses the distribution and allocation of collateral rather than its overall level.

2. These rules generally apply to what are known as QFCs, which include derivatives and repos. See Duffie and Skeel (2012), Skeel (2010), and Roe and Adams (2015) for legal background on stays for QFCs. See Bolton and Oehmke (2015), Duffie and Wang (2017) for corporate finance and game-theoretic models of automatic stays and contract termination.


4. This assumption is consistent with the treatment of IM in practice. The required IM is commonly held by a third-party custodian.

5. The assumption that collateral is taken first and the bankrupt firm’s assets are distributed in proportion to the residual claims is consistent with the legal discussion in Ayer et al. (2004).

6. Chang (2018) develops a model in which borrowers may fail to recover collateral when a lender defaults because the lender has, in turn, posted the collateral to another node. In our analysis, we assume nodes fully recover any collateral to which they are entitled.

7. It is also possible for node $i$ to default in round 1 yet end the round holding cash even before the return of collateral. However, in this case, we would have $\tilde{p}_{i}^{0} = 0$ for all $j$ and node $i$ has no remaining payments. This case is also illustrated in the figure.

8. Here and in several subsequent results, we are comparing two networks defined by fixed-point equations. To account for the possibility of multiple fixed points, the comparison should be understood to hold for the largest fixed points of the two networks and for the smallest fixed points. The existence of largest and smallest fixed...
points follows from applications of Tankiski’s (1955) fixed-point theorem. These issues are discussed in greater detail in the proofs of the relevant results.

8 Most OTC derivatives contracts are governed by an ISDA master agreement, which specifies the very limited settings—primarily events of default by the counterparty—in which one party may terminate a contract. See the discussion in Appendix B.

9 In formulating $V_1$, we have made the reasonable assumption that banks use their liquid assets before selling illiquid returned collateral. If the order is reversed, $V_1^+V_1^−$ becomes $V_1^+V_1^− sensual(V(V)_{sensual})$.

10 As explained in Skeel and Jackson (2012), debtors may engage in their own form of cherry-picking. With the protection of a bankruptcy stay, a debtor may decide which contracts to “assume” and which to “reject.” Skeel and Jackson (2012) also explain that this optionality is consistent with the treatment of “executory contracts” in nonfinancial bankruptcies.

11 Title II of the Dodd–Frank Act and the resolution framework of the Federal Deposit Insurance Corporation (in coordination with the Federal Reserve and Office of the Comptroller of the Currency (OCC)) impose a one- to two-day stay for QFC counterparties of the most complex U.S. bank holding companies (U.S. globally systemically important banks). Under the U.S. Treasury’s proposed chapter 14 bankruptcy process, termination rights of QFC counterparties are stayed for two days (U.S. Department of the Treasury 2018). That is, automatic stays for QFCs can be in place in both bankruptcy and resolutions proceedings.

12 Although we interpret selective termination as the result of creditor cherry-picking, a similar outcome would result if the failed node chose to terminate all its out-of-the-money contracts. This possibility is noted in Skeel and Jackson (2012). By terminating these contracts, a failed node could force its counterparties to accept lower payments through bankruptcy proceedings.

13 Our use of $q$ and $\pi$ in this section should not be confused with their use in equations 3.3. In both cases, we use $q$ and $\pi$ to indicate an alternative to a model that uses $p$ and $\pi$.

References


